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Performance Evaluation of WiMAX System Using Different Coding Techniques

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1. Introduction

In this chapter, we introduce a new class of coding technique that belongs to product code family. This technique is based on convolutional code. The use of convolutional code in the product code setting makes it possible to use the vast knowledge base for convolutional codes as well as their flexibility.

Product codes studied thus far have been constructed using linear block codes, such as Hamming [1], Bose–Chaudhuri–Hocquenghem (BCH) [2] and [3], Reed Solomon codes [4] and single parity check (SPC) [5]. These types of the product codes are traditionally constructed by linear block codes that have structure with a time varying property [6].

The product code proposed in this chapter is constructed by using time-invariant convolutional code. Its component codes' trellis structure does not vary in time as in product codes constructed with Hamming, BCH, and Reed Solomon block codes. Moreover, the number of states in the trellis structure of a block code may grow exponentially with the difference of codeword and data block lengths, whereas the number of states in a convolutional code can be set as desired.

The time invariant trellis structure of convolutional codes makes them more convenient for implementation. In addition, numerous practical techniques such as trellis coded modulation and puncturing can be simply utilized with convolutional codes as opposed to linear block codes.

Multi-input multi-output (MIMO) techniques are quite important to enhance the capacity of wireless communication systems. Space-time trellis codes provide both diversity and coding gain in MIMO channels and are widely used [7]. Space-time trellis codes usually have time-invariant trellis structures just like convolutional codes. Thus, a product code based on convolutional codes is more suitable for integration with MIMO channels and poses an alternative to block product codes.

The type of proposed product code described in this chapter is called modified Convolutional Product Codes (CPC), considered as a different type of normal CPC [8]. The normal CPC depends on recursive systematic convolutional encoder, whereas the modified version of CPC will basically depend on non-recursive non-systematic convolutional encoder.

WiMAX system is a wireless communication system. It suffers from having a high *Bit Error Rate* (BER) at low *Signal to Noise Ratio* (SNR). Using modified version of CPC for WiMAX system reduces the BER at low SNR. Also using the modified version of CPC with WiMAX decreases the number of stages of its physical layer as described later.

2. CPC encoder

For a regular product code, the information bits are placed into a matrix. The rows and columns are encoded separately using linear block codes. This type of a product encoder is shown in Figure 1.

Fig. 1. Regular Product Code Encoding Procedure.

In CPC method, the information bits are placed into two dimensions (2D) matrix. The rows and the columns are encoded separately using recursive systematic convolutional encoders. Each row of the matrix is encoded using a convolutional code with generator polynomial (1, 5/7) octal and code rate (1/2) Figure 2. The same recursive systematic convolutional code with the same polynomial is used to encode each row. Once all rows have been encoded, the matrix is sent, if desired, to an interleaver. The original data matrix dimensions are (n × k), and the encoded data matrix dimensions will be (2n×k) for coding rate (1/2).

The coded rows matrix is then recoded column by column using the same or different recursive systematic convolutional encoder. Hence, the overall code rate is 1/4.

Fig. 2. CPC Convolutional Coding [1 , 5/7].

The general encoding procedure, which includes any type of convolutional encoder and interleaver, is illustrated in Figures 3 & 4.

Fig. 3. CPC Encoding Procedure without an Interleaver.

Fig. 4. Convolutional Product Code Encoder with any type of Interleaver (d denotes data bits and p denotes parity bits).

3. CPC decoder

In the decoding process, the log-MAP soft decoding algorithm, [9] and [10], is used to iteratively decode the convolutional product code. Since columns were encoded last, each column is independently decoded one by one. The extrinsic information obtained from the columns is passed to the row decoder after being de-interleaved. Then, row decoding proceeds; rows are decoded one by one, and interleaved extrinsic information is passed to the column decoder. The CPC decoding procedure is depicted in Figure 5.

The decoding structure employed in this method is the same as that of serially-concatenated codes in Figure 6.6 [11].

4. Modified CPC encoder

In the modified version of CPC, the same technique is used for coding the message, except using nonrecursive nonsystematic convolutional encoder instead of recursive systematic

convolutional encoders for coding both rows and columns. That means the both encoders of rows and columns will have coding rate (1/2), and generator polynomial (5,7) Octa Figure 7.

Fig. 5. Decoding Operation of the Convolutional Product Code.

Fig. 6. Serial Encoding & Decoding Operations

The sequence of bits is fed into 2D matrix and fills it column by column. The size of this matrix depends only on the type of modulation used. For 16 QAM, the size of the matrix will be (nx4) and for 64 QAM the size of the matrix will be (nx6). These sizes simplify the process of mapping, as the symbol size in 16 QAM is 4 bits and in 64 QAM is 6 bits. So each row of those matrices will form one QAM symbol. The 'n' refers to the number of data subcarriers of OFDMA, 128 or 512.

The coding by modified CPC will be done in 2 stages

1. Each column will be independently coded.
2. Then each row of the resulting matrix will be coded by the same generator polynomials.

From Figure 7 since the generator polynomials used for coding both rows and columns are $(5,7)_{\text{octal}}$ with constraint length 3, not following the standard of WiMAX, each column is padded with two zeros for terminating its encoder. But each row is padded with two or three zeros according to the number of used subcarriers, 128 or 512, receptively to form the suitable size of the overall matrix. That matrix is then divided into smaller matrices with sizes $(n \times 4)$ or $(n \times 6)$ as described later.

Fig. 7. Convolutional Coding [5,7].

After the coding process, the total number of bits will be more than the original message bits due to the increase in the overall code rate $(1/4)$, and the addition of the zeros in both column and rows that used for the termination process. Therefore the following steps are done,

1. Dividing the overall matrix produced from modified CPC into three matrices. Each one has a size $(n \times 4)$ or $(n \times 6)$ according to the type of QAM used as mentioned before. The reason for using three matrices only is to have a number of message bits equals to bits used in the convolutional code method, as a comparison between it and CPC is done.
2. Applying symbol mapping for each one independently (16QAM or 64 QAM).
3. Inserting the pilot and DC subcarriers for each matrix.
4. Performing the IFFT on the three matrixes independently resulting in three OFDMA symbols.
5. Applying (cyclic prefix) CP for each symbol.
6. Sending each symbol independently.

The reason for using nonrecursive nonsystematic convolutional encoder instead of recursive systematic convolutional encoders is simplifying the termination of the encoder, as RSC contains a feedback and its termination will be more difficult. Also using the generator polynomials $(5,7)$ leads to a little increase in the complexity of the system because of a few number of zeros will be added to terminate the two encoders.

5. Modified CPC decoder

At the receiver, the three OFDMA symbols are combined to form the original matrix which is decoded by Viterbi decoder. The Viterbi decoder uses the same generator polynomials (5,7) with hard decision for each row and for each column. The rows must be decoded first then the columns are done, because columns are encoded first Figure 6. To match the CPC method, the number of data bits will be reduced. For example in OFDMA (128-16QAM) and (128-64QAM) the number of data bits was 144 and 216 but in CPC method it becomes 136 and 204 bits respectively due to the number of zero bits added to terminate the two encoders.

6. Modified CPC minimum distance and its asymptotic performance

The Hamming weight of a binary codeword is defined as the number of '1's available in the codeword. The minimum distance of a linear code is the minimum Hamming weight of all the codewords. The minimum distance plays an important role in the code performance. As it gets larger, code performance improves, especially at high signal-to-noise ratio (SNR) values. The free distance of the component convolutional codes used in modified CPC with trellis termination will be called d_{free} . The minimum distance of the modified CPC in the case of no interleavers will be investigated.

No Interleaving

After the first stage of the modified CPC encoding operation (columns encoding), it is obvious that one of these columns should contain at least d_{free} number of '1's. This means that there are d_{free} rows containing at least a *single* '1' in the columns-encoded matrix. When rows are encoded, there exists at least d_{free} number of rows each containing at least d_{free} '1's. Hence, in total there are at least d_{min}^2 '1's in the coded matrix. In summary, if no interleavers are used, the modified CPC minimum distance is d_{min}^2 .

7. Advantage and disadvantage of CPC

CPC technique has mainly two main advantages that make it a motivating step for future considerations and improvements for practical systems.

1. Do not need another interleaver after channel coding because of converting into matrix (nx4) or (nx6) does almost the same job as the overall matrix will be filled column by column and will be read row by row after coding processes (block interleaver) since each row is used for making QAM symbol.
2. Reducing the BER at low SNR.
3. The product code we propose in CPC is constructed by using time invariant convolutional codes. Its component codes' trellis structure does not vary in time as in product codes constructed with Hamming, extended Hamming, BCH, and Reed Solomon block codes. The time invariant trellis structure of convolutional codes makes them more convenient for implementation
4. The number of states in CPC like a convolutional code can be set as desired.
5. Numerous practical techniques such as trellis coded modulation and puncturing can be simply utilized with convolutional codes as opposed to linear block codes.

6. Space-time trellis codes usually have time-invariant trellis structures just like convolutional codes. Thus, a product code based on convolutional codes is more suitable for integration with MIMO channels.
7. Increasing the free distance to be $d_{\min}^2$.

But on the other hand it causes more delay for obtaining the original message because the code rate becomes 1/4 not 1/2 as in convolutional code. The performance of the system will be reduced and this is the price to be paid for the improvement obtained.

8. Results

This section contains comparisons between the modified CPC method and convolutional code, turbo code and LDPC code. Several results obtained at different types of the channels, modulation techniques (16QAM – 64QAM) and number of OFDM subcarriers (128 -512).

In this work, a matlab tool is used to simulate the physical layer of WiMAX and apply the mentioned coding methods.

8.1 AWGN channel

In this section the coded signal is transmitted through AWGN channel only. This can be done using matlab function **AWGN**. The syntax of this function is: $y = \text{awgn}(x, \text{snr}, 'measured')$ adds white Gaussian noise to the vector signal x to produce output signal y . The scalar snr specifies the signal-to-noise ratio per sample, in dB. If x is complex, **awgn** adds complex noise. This syntax measures the power of x before adding noise.

We can derive the relationship between E_s/N_0 and SNR for complex input signals as follows:

$$\begin{aligned}
 E_s / N_0 \text{ (dB)} &= 10 \log_{10} ((S \cdot T_{\text{sym}}) / (N / B_n)) \\
 &= 10 \log_{10} ((T_{\text{sym}} F_s) \cdot (S / N)) \\
 &= 10 \log_{10} (T_{\text{sym}} / T_{\text{samp}}) + \text{SNR (dB)}
 \end{aligned} \tag{1}$$

Where

S = Input signal power, in watts

N = Noise power, in watts

B_n = Noise bandwidth, in Hertz

F_s = Sampling frequency, in Hertz = $1/T_{\text{samp}}$.

T_{samp} = The period of each row of a frame-based matrix.

T_{sym} = The signal's symbol period.

A good rule of thumb for selecting the symbol period value is to set it to be what we model as the symbol period in the model. The value would depend upon what constitutes a symbol and what the oversampling applied to. From Figure 8 to Figure 13 BER versus different received SNR values are shown for the comparison between modified CPC and

convolutional code, LDPC code and turbo code respectively. These comparisons are obtained for modulation type 16QAM and number of subcarriers equals 128 and 512 respectively.

The comparisons between modified CPC and convolutional code are shown in both Figure 8 and Figure 9. From Figure 8, it is shown that SNR will be improved by approximately 2 dB at BER equals 10^{-3} for modulation type 16QAM and number of subcarriers equals 128. Also, an improvement can be obtained when the number of subcarriers is increased to 512 as shown in Figure 9.

Fig. 8. BER Comparison between Conv code, CPC at 16 QAM, N=128.

Fig. 9. BER Comparison between Conv code, CPC at 16QAM, N=512.

From Figure 10 and Figure 11, the results of comparisons between modified CPC and LDPC code are shown at different received SNR values. From these figures, we conclude that modified CPC gives good results at different SNR. Figure 10 shows that modified CPC coding technique gives better results than LDPC coding technique at 16 QAM and OFDM subcarriers equals 128. The improvement is more than 3 dB for 10^{-3} . Also, there will be an improvement, when the number of subcarriers is increased to 512 as shown in Figure 11.

Fig. 10. BER Comparison between LDPC code, CPC at 16 QAM, N=128.

Fig. 11. BER Comparison between LDPC code, CPC at 16 QAM, N=512

From Figure 12 to Figure 13, the results produced from the comparisons between modified CPC and turbo code are shown at different received SNR values. As shown from these figures, modified CPC method gives good results compared to turbo coding. Figure 12 shows that using modified CPC method can give better results than turbo coding technique at 16QAM and OFDM subcarriers equals 128. This improvement is more than 3 dB for BER= 10^{-3} . Also other improvements can be obtained at different number of OFDM subcarriers (512) as shown in Figure 13.

Fig. 12. BER Comparison between Turbo code, CPC at 16 QAM, N=128.

Fig. 13. BER Comparison between Turbo code, CPC 16 QAM, N=512.

The comparison for modulation type 64 QAM between modified CPC and convolutional code, LDPC code and turbo code are shown through Figure 14 to Figure 19. BER versus different received SNR values are shown in these figures. These comparisons are obtained and number of subcarriers equals 128 and 512 respectively. The comparisons between modified CPC and convolutional code are shown in both Figure 14 and Figure 15.

Fig. 14. BER Comparison between Conv code, CPC at 64 QAM, N=128.

Fig. 15. BER Comparison between Conv code, CPC at 64 QAM, N=512.

From Figure 14, it is shown that SNR will be improved by approximately 1.5 dB at BER equals 10^{-2} for modulation type 64QAM and number of subcarriers equals 128. Also, an

improvement can be obtained when the number of subcarriers is increased to 512 as shown in Figure 15. Figure 16 and Figure 17 show the results of comparisons between modified CPC and LDPC code are shown at different received SNR values. We conclude that modified CPC gives good results at different SNR. Figure 16 shows that modified CPC coding technique gives better results than LDPC coding technique at 64QAM and OFDM subcarriers equals 128. The improvement is more than 1.5 dB for 10^{-2} . Also, there will be an improvement, when the number of subcarriers is increased to 512 as shown in Figure 17.

Fig. 16. BER Comparison between LDPC code, CPC at 64QAM, N=128.

Fig. 17. BER Comparison between LDPC code, CPC at 64 QAM, N=512.

From Figure 18 to Figure 19, the results produced from the comparisons between modified CPC and turbo code are shown at different received SNR values. As shown from these figures, modified CPC method gives good results compared to turbo coding. Figure 18 shows that using modified CPC method can give better results than turbo coding technique at 64QAM and OFDM subcarriers equals 128. This improvement is about than 2 dB for 10^{-2} . Also other improvements can be obtained at different number of OFDM subcarriers (512) as shown in Figure 19.

Fig. 18. BER Comparison between Turbo code, CPC at 64 QAM, N=128.

Fig. 19. BER Comparison between Turbo code, CPC 64 QAM, N=512.

8.2 AWGN plus fading channel

In this section, the transmitted signal is assumed to pass through time selectivity fading channel plus AWGN. This is done using matlab function **rayleighchan**. The syntax of this function is **chan = rayleighchan (Ts,Fd)** that constructs a frequency-flat ("single path") Rayleigh fading channel object.

Ts is the sample time of the input signal, in seconds. Fd is the maximum Doppler shift, in Hertz.

$$\text{Sample time} = 1/(\text{channel bandwidth} \times 28/25) \quad (2)$$

$$\text{Maximum Doppler shift} = F_d = (v f_c) / C_0 \quad (3)$$

Where f_c is carrier frequency, v is the maximum speed between transmitter and the receiver and C_0 is the speed of light. The Rayleigh multipath fading channel simulators of this toolbox use the band-limited discrete multipath channel model. It is assumed that the delay power profile and the Doppler spectrum of the channel are separable. The multipath fading channel is therefore modeled as a linear finite impulse-response (FIR) filter. Let S_i denotes the set of samples at the input to the channel. Then the samples y_i at the output of the channel are related to S_i through:

$$y_i = \sum_{n=-N_1}^{N_2} s_{i-n} g_n \quad (4)$$

Where g_n is the set of tap weights given by:

$$g_n = \sum_{k=1}^K a_k \text{sinc} \left[\frac{\tau_k}{T_s} - n \right], \quad -N_1 \leq n \leq N_2 \quad (5)$$

In the equations above:

- T_s is the input sample period to the channel.
- τ_k , where $1 \leq k \leq K$, is the set of path delays. K is the total number of paths in the multipath fading channel.
- a_k , where $1 \leq k \leq K$, is the set of complex path gains of the multipath fading channel. These path gains are uncorrelated with each other.
- N_1 and N_2 are chosen so that $|g_n|$ is small when n is less than N_1 or greater than N_2

This simulation is done for different coding techniques that have different coding rates because we follow the standard in our simulation. The following parameters are used in our simulation:

1. Frequency band is 3.5 GHz.
2. Channel Bandwidth (1.25 MHz for IFFT size=128 and 5.00 MHz for IFFT size= 512).
3. Modulation types (16 QAM, 64 QAM).
4. Oversampling rate is 28/25.
5. Max speed 120 Kmph.
6. Convolutional code with rate equals (1/2), turbo code with rate equals (2/3) and LDPC code with rate equals (1/2).

From Figure 20 to Figure 25 BER versus different received SNR values are shown for the comparison between modified CPC and convolutional code, LDPC code and turbo code respectively through the fading channel. These comparisons are obtained for modulation type 16QAM and number of subcarriers equals 128 and 512 respectively.

The comparisons between modified CPC and convolutional code are shown in both Figures 20 and 21. In Figure 20, it is shown that SNR is improved by more than 4 dB at BER equals 10^{-2} for the number of subcarriers equals 128. An improvement is obtained if the number of subcarriers is increased to 512 as shown in Figure 21.

Fig. 20. BER Comparison between Conv code, CPC at 16QAM , N=128.

Fig. 21. BER Comparison between Conv code, CPC at 16 QAM, N=512.

The results of comparisons between CPC and LDPC code through the fading channel are obtained from Figure 22 to Figure 23. These comparisons are obtained for modulation 16QAQM at different SNR values. From Figure 22, it is shown that SNR is improved by about 2.5 dB at BER equals 10^{-2} for the number of subcarriers equals 128. Another improvement is also obtained at different number of OFDM subcarriers (512) as shown in Figure 23.

Fig. 22. BER Comparison between LDPC code, CPC at 16 QAM, N=128.

Fig. 23. BER Comparison between LDPC code, CPC at 16 QAM, N=512.

The results of comparisons between modified CPC and turbo code through the fading channel are shown from Figure 24 to Figure 25. These comparisons are done at different

SNR values for modulation type 16QAM. There is an improvement in SNR by more than 8 dB at BER equals 10^{-2} for 16QAM and number of OFDMA subcarriers equals 128, this is shown from Figure 24. Other improvements obtained at different number of OFDMA subcarriers (512) as shown from Figure 25.

Fig. 24. BER Comparison between Turbo code, CPC at 16QAM , N=128.

Fig. 25. BER Comparison between Turbo code, CPC at 16 QAM, N=512.

The comparison for modulation type 64QAM between modified CPC and convolutional code, LDPC code and turbo code are shown through Figure 26 to Figure 31 through the

fading channel.. BER versus different received SNR values are shown in these figures. These comparisons are obtained and number of subcarriers equals 128 and 512 respectively. The comparisons between modified CPC and convolutional code through the fading channel are shown in both Figure 26 and Figure 27.

Fig. 26. BER Comparison between Conv code, CPC at 64QAM, N=128.

Fig. 27. BER Comparison between Conv. code, CPC at 64 QAM, N=512.

From Figure 26, it is shown that SNR will be improved by more than 2 dB at BER equals 10^{-2} for modulation type 64QAM and number of subcarriers equals 128. Also, an improvement

can be obtained when the number of subcarriers is increased to 512 as shown in Figure 27. Figure 28 and Figure 29 show the results of comparisons between modified CPC and LDPC code through the fading channel are shown at different received SNR values. We conclude that modified CPC gives good results at different SNR. Figure 28 shows that modified CPC coding technique gives better results than LDPC coding technique at 64QAM and OFDM subcarriers equal 128. The improvement is more than 1.5 dB for 10^{-2} . Also, there will be an improvement, when the number of subcarriers is increased to 512 as shown in Figure 29.

Fig. 28. BER Comparison between LDPC code, CPC at 64 QAM, N=128.

Fig. 29. BER Comparison between LDPC code, CPC at 64 QAM, N=512.

From Figure 30 to Figure 31, the results produced from the comparisons between modified CPC and turbo code are shown at different received SNR values for modulation 64 QAM. As shown from these figures, modified CPC method gives good results compared to turbo coding. Figure 30 shows that using modified CPC method can give better results than turbo coding technique at 64QAM and OFDM subcarriers equals 128. This improvement is about 2 dB for 10^{-2} . Also other improvements can be obtained at different number of OFDM subcarriers (512) as shown in Figure 31.

Fig. 30. BER Comparison between Turbo code, CPC at 64 QAM, N=128.

Fig. 31. BER Comparison between Turbo code, CPC at 64 QAM, N=512.

Due to the lower code rate of CPC (1/4), better results should be obtained comparing to the other coding techniques (Convolution – Turbo – LDPC) which has higher coding rate (1/2 or 1/3). But from the view of complexity, CPC is less complex than turbo or LDPC, so the type of coding can be used as an optional code instead of turbo or LDPC types. Another technique can be used with modified CPC to modify or increase its code rate, this technique called *puncture technique*.

9. Punctured CPC

To modify the rate of the coding process, a puncture technique is used. This technique enables to have a code rate equals 1/3. The modification is done by applying the puncture to the columns only, resulting in code rate of 2/3. So the overall code rate will be (2/3) x (1/2) = (1/3). Puncture enables to reduce the redundancy bits but on other hand it leads to increase the BER. From Figure 32 to Figure 42 the result of using CPC with puncture is shown, through AWGN plus fading channel, comparing with convolutional code, turbo code and LDPC code. From Figure 32 to Figure 34 BER versus different received SNR values are shown for the comparison between modified CPC, puncture CPC and convolutional code, LDPC code and turbo code respectively through the fading channel. These comparisons are obtained for modulation type 16QAM and number of subcarriers equals 128.

Fig. 32. BER Comparison between Conv code, CPC, punctured CPC at 16 QAM, N=128.

Fig. 33. BER Comparison between LDPC code, CPC, punctured CPC at 16 QAM, N=128.

Fig. 34. BER Comparison between Turbo code, CPC, punctured CPC at 16 QAM, N=128.

From Figure 32 it is shown that the results obtained from LDPC code is approximately the same as puncture modified CPC, but LDPC is more complicated than the proposed method. From Figure 35 to Figure 37 BER versus different received SNR values are shown for the comparison between modified CPC, puncture CPC and convolutional code, LDPC code and turbo code respectively through the fading channel. These comparisons are obtained for modulation type 16QAM and number of subcarriers equals 512.

Fig. 35. BER Comparison between Conv code, CPC, punctured CPC at 16 QAM, N=512.

Fig. 36. BER Comparison between LDPC code, CPC, punctured CPC at 16 QAM, N=512.

Fig. 37. BER Comparison between Turbo code, CPC, punctured CPC at 16 QAM, N=512

From Figure 36 it is shown that the results obtained from LDPC code is better than puncture modified CPC until SNR = 12 db, but LDPC is more complicated than the proposed method. From Figure 35 to Figure 37 BER versus different received SNR values are shown for the comparison between modified CPC, puncture CPC and convolutional code, LDPC code and turbo code respectively through the fading channel. These comparisons are obtained for modulation type 64QAM and number of subcarriers equals 128.

Fig. 38. BER Comparison between Conv code, CPC, punctured CPC at 64 QAM, N=128.

Fig. 39. BER Comparison between LDPC code, CPC, punctured CPC at 64 QAM, N=128.

From Figure 39 it is shown that the results obtained from LDPC code is better than t puncture modified CPC until SNR = 16db, but LDPC is more complicated than the proposed method. From Figure 35 to Figure 37 BER versus different received SNR values are shown for the comparison between modified CPC, puncture CPC and convolutional code, LDPC code and turbo code respectively through the fading channel. These comparisons are obtained for modulation type 64QAM and number of subcarriers equals 512.

Fig. 40. BER Comparison between Turbo code, CPC, punctured CPC at 64 QAM, N=128.

Fig. 41. BER Comparison between Conv code, CPC, punctured CPC at 64 QAM, N=512.

Fig. 42. BER Comparison between LDPC code, CPC, punctured CPC at 64 QAM, N=512.

From Figure 42 and Figure 43 it is shown that the results obtained from LDPC code and turbo code is better than puncture modified CPC, but LDPC and turbo code is more complicated than the proposed method.

Fig. 43. BER Comparison between Turbo code, CPC, punctured CPC at 64 QAM, N=512.

10. Delay diversity scheme

In this section, both puncture and Delay Diversity Scheme (DDS) are together used to increase the efficiency of CPC system. To further improve the diversity of the channels a transmit diversity technique may be utilized.

Many transmit diversity techniques have been explored. One such technique is the transmit delay diversity Figure 6.44. In transmit delay diversity a transmitter utilizes two antennas that transmit the same signal, with the second antenna transmitting a delayed replica of that transmitted by the first antenna. By so doing, the second antenna creates diversity by establishing a second set of independent multipath elements that may be collected at the receiver.

Fig. 44. Transmit Delay Diversity Scheme

If the multipath generated by the first transmitter fades, the multipath generated by the second transmitter may not, in which case an acceptable SNR will be maintained at the receiver. This technique is easy to implement, because only the composite TX0+TX1 channel is estimated at the receiver. The biggest drawback to transmit delay diversity is that it increases the effective delay spread of the channel, and can perform poorly when the multipath introduced by the second antenna falls upon, and interacts destructively with, the multipath of the first antenna, thereby reducing the overall level of diversity.

Our simulator for delay diversity technique is based on passing the same signal through the same path during two time intervals by using only one transmitted antenna Figure 6.45, not two transmitted antennas as in delay diversity technique. During these intervals the channel will have different fading and AWGN characteristics over the time.

Fig. 45. Modified Transmit Delay Diversity Scheme

The receiver, by using one received antenna, chooses the best received signal according to its highest power.

11. Conclusions

In this chapter, we explained CPC method as a coding technique and our modification for it. Also the implementation of CPC in WiMAX system and the comparisons between its results and the results of other coding techniques such as convolutional, turbo and LDPC are investigated at different SNR for different number of subcarriers and at different types of modulation (16QAM – 64QAM).

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Advanced Transmission Techniques in WiMAX

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This book has been prepared to present the state of the art on WiMAX Technology. The focus of the book is the physical layer, and it collects the contributions of many important researchers around the world. So many different works on WiMAX show the great worldwide importance of WiMAX as a wireless broadband access technology. This book is intended for readers interested in the transmission process under WiMAX. All chapters include both theoretical and technical information, which provides an in-depth review of the most recent advances in the field, for engineers and researchers, and other readers interested in WiMAX.

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