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Content Defined Optical Network

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<http://dx.doi.org/10.5772/intechopen.72432>

Abstract

Optical interconnection has become one of the key technologies to adapt the needs of large-scale data center networking with the advantages of large capacity, high bandwidth, and high efficiency. Data center optical interconnection has the characteristics of resource and technology heterogeneity. Its networking and control face enormous challenges for the increasing number of users with a high level quality of service requirements. Around different scenarios, there are a series of key networking and control problems in data center optical interconnection, such as multiple layers and stratum resources optimization in inter-data center, and time-aware resource scheduling in intra-data center. To solve these problems and challenges, this chapter mainly researches on content defined optical networking and integrated control for data center. For networking of vertical “multi-layer-carried” and horizontal “heterogeneous-cross-stratum”, the chapter launches research work around application scenarios about inter-data center optical interconnection with optical network, and intra-data center. The model architecture, implementation mechanism and control strategy are analyzed and demonstrated on the experiment and simulation platform of data center optical interconnection. This chapter will provide important references for future diverse applications of data center optical interconnection and software defined networking and control in practice.

Keywords: software defined optical network, content, data center, optical interconnect, OpenFlow

1. Introduction

With the rapid development of cloud computing and high rate services, data center services have attracted a great deal of attention from network service providers. With the variety and massiveness of applications, the high-performance network-based datacenter applications have the features of high burstiness and wide-bandwidth, particularly for the super-wavelength services [1]. Flexi-grid optical networking with big capacity, low power density and

distance adaption, provides a promising solution for the new datacenter network bottlenecks. A novel optical orthogonal frequency division multiplexing (OFDM)-based architecture with high spectral efficiency and high energy efficiency is presented for data center networks [2]. Networking architecture, algorithm and control plane for inter-datacenter network is also addressed in flexi-grid optical networks [3]. Lower power, improved scalability and port density which is the advantages of software defined optical networking for highly virtualized datacenters are studied [4].

Compared with optical interconnect networks between datacenters, optical interconnect networks in a datacenter is a more imperative requirement and respective case to serve the services in a flexible and high-efficient way [5]. Besides, miscellaneous datacenter services have a lower delay and higher availability requirements whose quality of service (QoS) can be guaranteed in corresponding levels [6]. Many studies focus on the architecture and equipment in datacenter interconnection [7–9]. For instance in Ref. [7], arrayed waveguide grating router (AWGR)-based interconnect architecture is proposed. A distributed all-optical control plane is designed with low latency and high-throughput at high traffic load in the case of sufficient packet transmission time. The authors in [8] achieve the hitless adaptation between Ethernet and time shared optical network (TSON) by designing a novel network on-and-off chip approach for highly efficient and transparent intra-datacenter communications. The work [9] make the datacenter offloads heavy inter-pod traffic onto an optical multi-ring burst network by proposing efficient scheme to all-optically inter-networking the pods. However, the time feature of application to guarantee the services delivery with various QoS in intra-datacenter networks from the view of service is relatively unexplored. Recently, as a centralized software control architecture, the software defined networking (SDN) enabled by OpenFlow protocol has become a focus of study by making the network functions and protocols programmable [10–13], where maximum flexibility is provided for the network operators and the integrated optimization of services can be achieved in a centralized control over multi-dimensional resource [14–17]. Hence, to introduce SDN method to centrally control network and application resources in optical interconnect of intra-datacenter has a great significance.

In our previous study, the network architecture with cross stratum optimization (CSO) based on SDN including multi-stratums resources in inter-datacenter networks has been designed to partially satisfy the QoS requirement [18–21]. Even in the edge of the network, the similar architectures with CSO based on SDN have been studied for improving the performance of cloud-based radio access network, which is similar as the inter-datacenter networks [22–24]. Based on the previous work, this paper proposes a Content Defined Optical Network (CDON) architecture in OpenFlow-based datacenter optical networks for service migration, in which a time-aware service scheduling (TaSS) strategy is introduced. CDON considers the time factor, in which the applications with required QoS can be arranged and accommodated to enhance the responsiveness to quickly provide for datacenter demand. By the experimental implementation on our testbed with OpenFlow-based intra-datacenter and inter-datacenter optical networks and the statistics collection of blocking probability and resource occupation rate, the overall feasibility and efficiency of the proposed architecture are verified. Intra-datacenter and inter-datacenter networks are considered in this paper. Based on the unified and flexible control advantages, SDN is deployed in both two networks.

The rest of this paper is organized as follows. In Section II, we propose the novel data-center-network-based CDON and builds functional models. Then intra-datacenter optical interconnection architecture and inter-datacenter optical network architecture are designed. Section III describe the TaSS strategy. Finally, we describe the testbed and present the experimental results and analysis in section IV and section V conclude the paper.

2. Datacenter-network-CDON

In order to promote the control efficiency of datacenter networks, control architecture based on CDON is described as shown in **Figure 1**. Different OpenFlow controllers for different resources including intra-datacenter computing recourse and inter-datacenter communication resources have been developed. The latter is mainly flexi-grid optical network resource in this work. All resources are software-defined with OpenFlow and support datacenter application. Then OpenFlow-enabled network controller and OpenFlow-enabled application controller can work together. By using user and application interface (UAI), application plane which is served through application and controller interface (ANI) can provide users with various services. The architecture of intra-datacenter and inter-datacenter networks is discussed in detail as follows.

2.1. Intra-datacenter optical interconnection architecture

The CDON architecture for OpenFlow-based intra-datacenter optical interconnect is shown in **Figure 2(a)**. Top-of-rack (ToR), aggregation and core optical switches three kinds of optical switches are used to interconnect datacenter servers with the deployment of application stratum resources (e.g., CPU and storage). Application controller (AC) and network controller (NC) respectively centrally control each stratum resources which are software defined with OpenFlow. To control intra-datacenter networks for service migration with

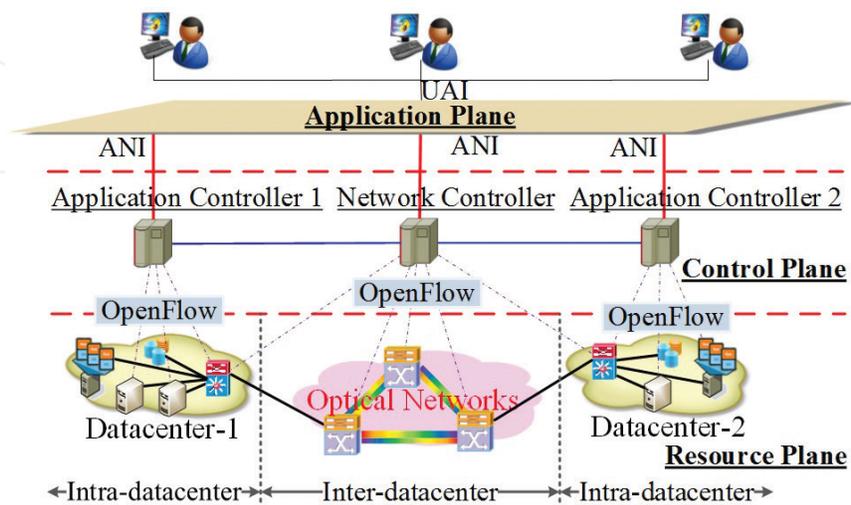


Figure 1. Datacentre-network-based CDON.

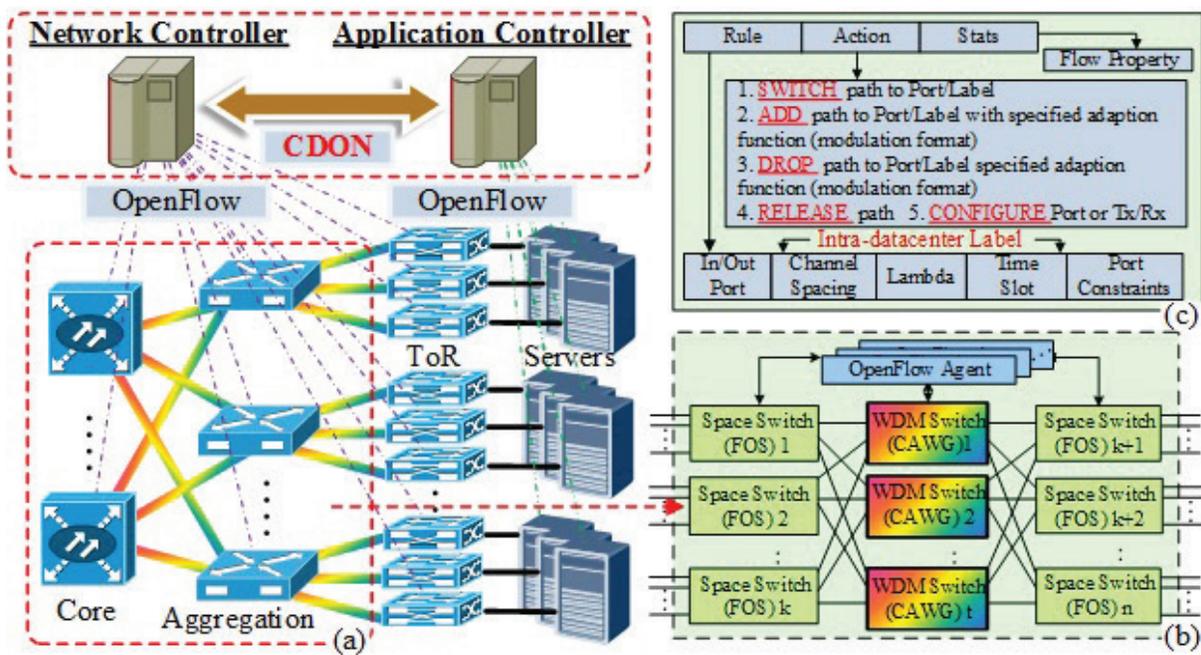


Figure 2. (a) The SDN-based intra-datacenter network architecture of CDON. (b) The SDN-based switching fabric of intra-datacenter. (c) The extension of protocol for CDON.

extended OpenFlow protocol (OFF), OpenFlow-enabled optical switches with the OFF agent software (OF-OS) is used. The CDON architecture in intra-datacenter networks have twofold motivations. Firstly, CDON can highlight the cooperation between AC and NC for supporting TaSS strategy to schedule datacenter applications based on different time sensitivity requirements reasonably and optimize application and network stratum resources efficiency. Secondly, considering the burstiness, burst applications fast provisioning with unified CDON control and process can be supported.

Figure 2(b) describes the OpenFlow-switching structure of intra-datacenter. The functions and interaction descriptions of relevant functional modules are described below. The AC is responsible for monitoring and maintaining the resources of application stratum for CDON, NC supports abstraction of the network information from physical stratum and lightpath provisioning of optical networks in intra-datacenter. When request arrives, AC processes it by using TaSS strategy considering the requirements of delay sensitivity and achieving CSO of the computing and storage in an internal database and sent the decision to NC. The control of virtual and physical network are concerned in NC. The former controls the virtual network and send abstracted network information to AC, the responsibilities of latter is to monitor and control the programmable physical modules. With the extended OFF, the lightpath can be built according to the request from AC. It is noteworthy that burst traffic can be correspondingly serviced by the high level optical switch by fast tunable laser (FTL) and burst mode receiver (BMR) in ToR switch. The aggregation and core switching fabrics are built based on space and wavelength circuit-switching technologies, fast optical switch (FOS) and cyclic arrayed waveguide grating (CAWG) respectively. The OFF agent software embedded in optical module has four functions which are maintaining flow table, modeling the information of node with programmability, mapping the content to configure and controlling hardware.

In connection to networks control of intra-datacenter, **Figure 2(c)** shows extended flow entry of OFP. The rule is added with the main characteristics of intra-datacenter including the in/out port, intra-datacenter label (e.g., channel space, lambda and time slot) and port constraints. The action is extended as five types: add, switch, drop and configure to set up a path, and release a path. Using combinations of rule and action, the control of optical node is realized. The responsibility of stats function is monitoring the flow property to provide service provisioning for CDON.

2.2. Optical network architecture of inter-datacenter

Flexi-grid optical networks are the promising technology for inter-datacenter networks as these networks can satisfy the requirement of burstiness. The CDON architecture is built and shown in **Figure 3(a)**. The distributed datacenters are interconnected through the flexi-grid optical networks. The network architecture mainly consists of the optical resources stratum and application resources stratum. In a unified manner, network controller and application controller can control each resource stratum which is software defined with OpenFlow. Software defined OTN (SD-OTN) is necessary to control the flexi-grid optical networks for inter-datacenter network with extended OFP. SD-OTNs are essentially OpenFlow-enabled elastic optical device nodes with OFP agent software. It has twofold motivations to design the CDON architecture over inter-datacenter optical network. Firstly, the CDON can realize the global interworking of cross stratum resources that the physical

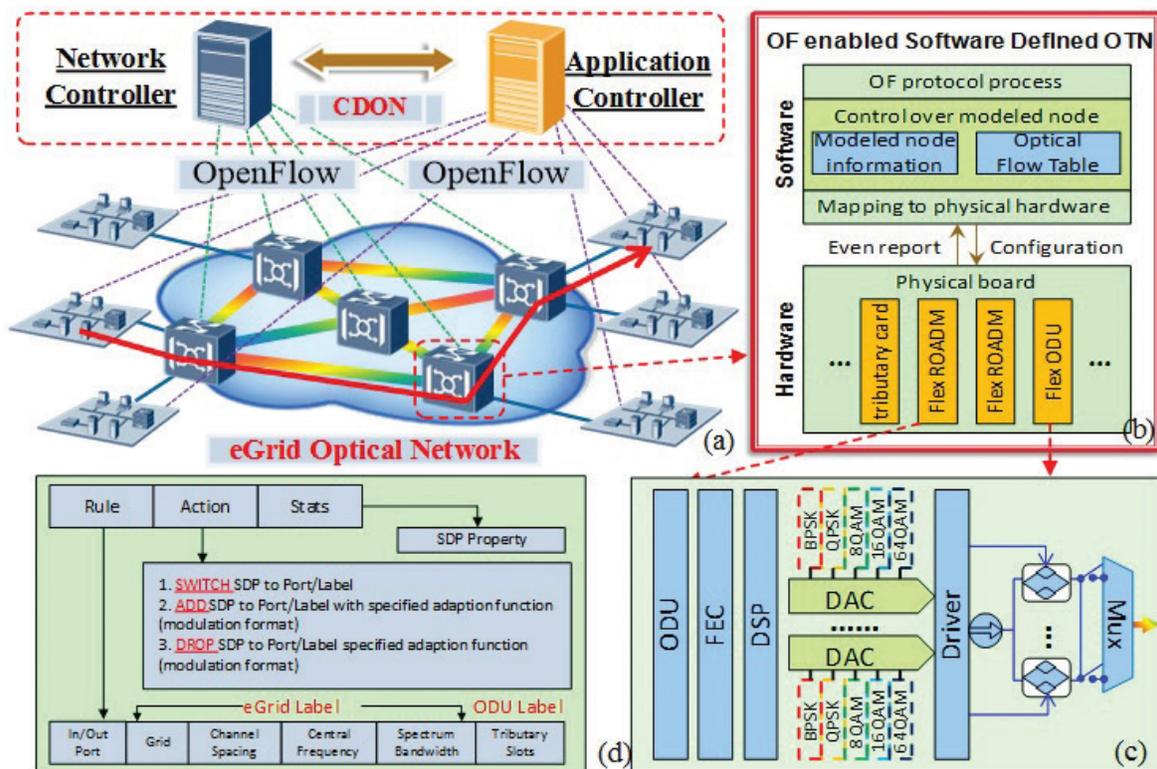


Figure 3. (a) The architecture of CDON in OpenFlow-based inter-datacenter networks. (b) OF-enabled SD-OTN functional models. (c) Flex ODU. (d) The extension of protocol for CDON in inter-datacenter.

layer parameter (e.g., bandwidth and modulation format) can be adjusted. So the cooperation between AC and NC is emphasized to realize software defined path (SDP) with application and spectrum elasticity. Secondly, the different time sensitivity requirements of services can be considered reasonably through scheduling data center services with time elasticity to optimize the application and network resources utilization further. Based on functional architecture described above, TaSS scheme is proposed in the AC and it can arranging the start time, transport time and corresponding transport bandwidth for services for realizing the application and network stratum resources optimization.

The functional modules of AC and NC and the coupling relationship between different modules are shown as follow. NC is responsible for the control in physical and virtual network. The former includes controlling spectrum resource and modulation format. The latter is responsible for managing the virtual network and sending virtual resource information to AC. While a request arrives, AC runs TaSS strategy based on multi-stratum resource information and sent the decision to NC through application-transport interface (ATI). The SDP can be found out and built up based on extended OFP according to the request from AC. It is noteworthy that, the length of SDP decides the modulation format of service (e.g., QPSK and 16QAM). In the case of short distance, spectrum bandwidth which is more precious than other resource can be economized by using high-level modulation format. In SD-OTN, OpenFlow-enabled agent software is embedded to realize the communication between NC and optical node. The SD-OTN maintains optical flow table and modeled node information as software. The physical hardware which contains flexible ROADMs and ODU boards shown in **Figure 3(b)** and **(c)** respectively is configured and controlled through the content mapping. In the side of the control of flexi-grid optical networks in inter-datacenter, flow entry of OFP is extended and shown in **Figure 3(d)**. In this architecture, the rule is extended as the main characteristics of flexi-grid optical networks which including the in/out port, flexi-grid label (e.g., central frequency and spectrum bandwidth) and ODU label (e.g., tributary slots). The action of optical node mainly includes three types: add, switch and drop. Through combining rule and action, the control of flexi-grid node can be realized. The responsibility of state function is monitoring the flow property to provide SDP provisioning.

3. Service scheduling strategy

In the side of the service accommodation of datacenter optical networks, the traditional strategy can.

allocate the optimal datacenter server application and corresponding lightpath network resources when a request arrives. That may satisfy the following conditions, which are shown in **Figures 4** and **5**. We assume that two datacenter servers are deployed in the candidate destination node. The storage utilization of server #1 and server #2 are 80% and 85% respectively when service arrives. After a relatively short time, the services which are provided by the two servers have changed, i.e., some services are released when they are complete and new ones arrive. It leads that the storage utilization of server #1 increases to 95%, while that in server #2 is down to 25%. The server #1 (80% at arriving time) would be chosen as the destination

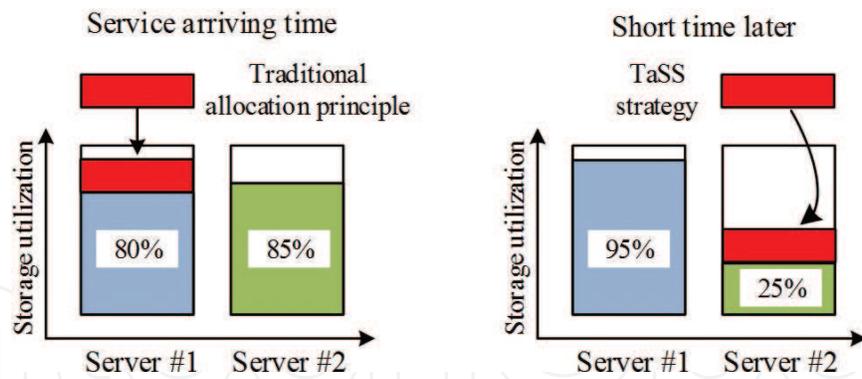


Figure 4. Illustration of time-aware datacenter application resource allocation.

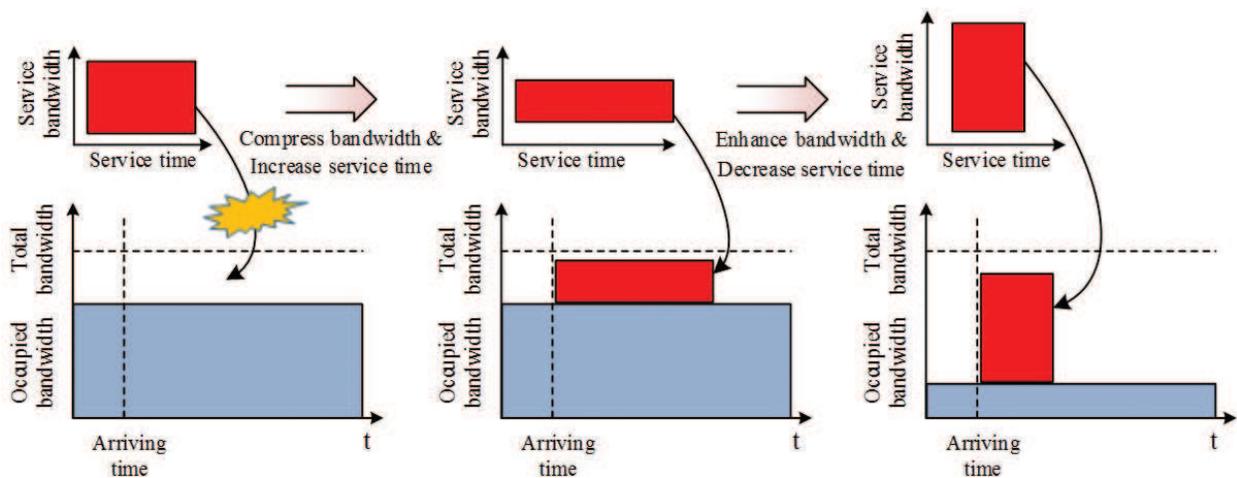


Figure 5. Illustration of time-aware network resource allocation.

node When the traditional allocation principle is used. However, if the resource is allocated after a short time (within delay tolerance of user), the better optimization and more effective resource utilization may be realized. The other instance is about the service bandwidth. The bandwidth which most of datacenter services use is specified and the time to transmit service is fixed for achieving the overall data volume of service. If the bandwidth is fixed with traditional methods, the service cannot be provided with enough resource in case of less available bandwidth. With compressing the bandwidth to adapt the available bandwidth and increasing the service time (overall data volume of service is constant), the allocation scheme can be feasible. Similarly, by enhancing the provide bandwidth, the service can be completed as soon as possible with relatively abundant resources. The time factor of datacenter service is considered in the two issues. Therefore, we propose the time-aware service scheduling strategy.

3.1. Network modeling

$G(V, L, F, A)$ denotes the OpenFlow-based datacenter interconnect with optical networks, where $V = \{v_1, v_2, \dots, v_n\}$ is represented as the set of SDN-based optical switching nodes, $L = \{l_1, l_2, \dots, l_n\}$ denotes the set of optical fiber links connecting nodes in V . $F = \{\omega_1, \omega_2, \dots, \omega_F\}$

indicate the set of wavelengths of optical fiber and A is the set of servers in datacenter. While source node s sends a service request, the request contains total data volume of service D and storage space S . The services are classified with latency-sensitive and delay-tolerant service. The former requires immediate service process and its i th request is denoted as $SR_i(s, D, S)$, and the latter includes the arriving time t_c and tolerant delay T , and we denote its i th service request as $SR_i(s, D, S, t_c, T)$. The request SR_{i+1} will be the next request in time order when the connection demand SR_i arrives. In addition, **Table 1** shows some requisite notations and their definitions.

3.2. Time-aware service schedule strategy

Developing on the functional architecture, we present a novel time-aware service scheduling (TaSS) strategy which is implemented in application controller to schedule datacenter service with time sensitivity requirement. For the requirements of the arriving service, we classify them as burst latency-sensitive and delay-tolerant service including flow volume and tolerant latency, based on the delay sensitivity of each service. For the delay-tolerant service, we search servers and lightpaths of datacenter to judge whether the volume of S_r , B_r and t is enough to satisfy the incoming service. If the $S_r \geq S$ and $B_r \cdot t \geq D$, by comparing CSO factor [16], the minimum one would be chosen to provision. If the resources for accommodating is not enough, the service would wait until maximum tolerant delay arrives, i.e., $t = T - D/B_a - t_c$, which considers B_a and t_c . New lightpath will be allocated with CSO to accommodate the service. In the side of delay-sensitive requests, by searching the existing candidate with the same method, the available resources are decided. The path which has minimum propagation delay obtained from candidate paths can be prepared to provision. When an available built path is found, the TaSS.

Strategy builds new lightpath for the service immediately. In order to realize the quick response for service provisioning, service time is used to reckon the release lightpath procedure immediately. The flowchart of TaSS strategy is shown in **Figure 6**.

Symbol	Definitions
S	The total capacity of storage space in target data center
S_r	The residual storage space in target data center
D	The total data volume of service
B_r	The product of available bandwidth of lightpath
T	The tolerant delay of delay sensitive service
t	The lightpath duration
B_a	The network average transmission bandwidth
t_c	The guard time
N	The number of network nodes
L	The number of links
F	The number of wavelengths
A	The number of datacenter nodes

Table 1. Symbols and definitions.

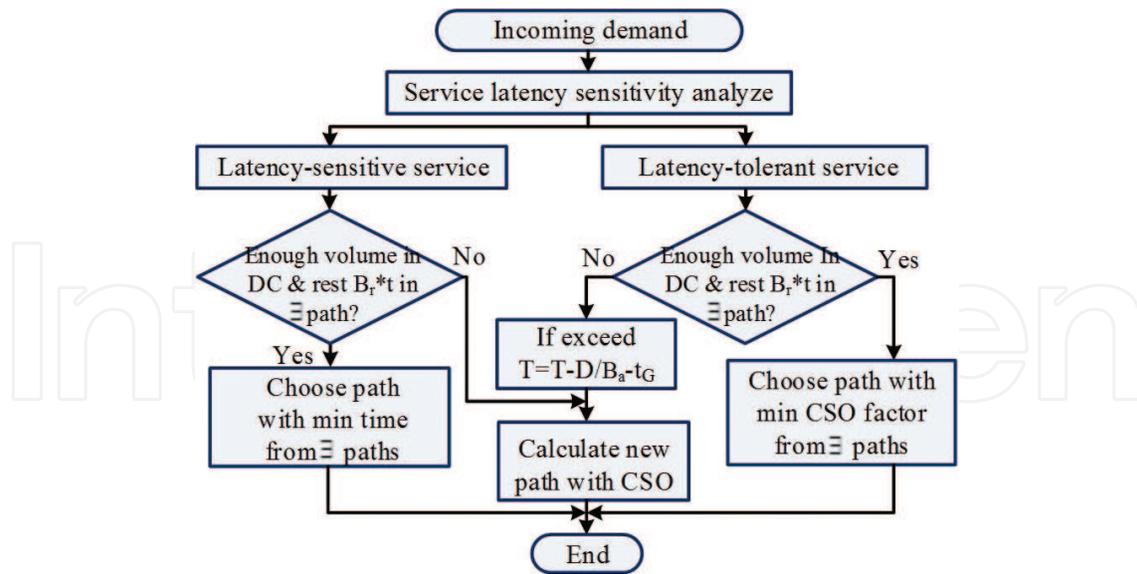


Figure 6. Flowchart of TaSS strategy.

4. Experimental demonstration and results discussion

A testbed of datacenter networks consisting of intra-datacenter and inter-datacenter is built. SDN is deployed on the two parts together. The following is a detailed description of the testbed demonstrations.

For experimentally evaluating CDON architecture, we deploy it into optical intra-datacenter networks and realize the service migration in this architecture on our testbed as shown in **Figure 7(a)**. In data plane, it contains 4 optical switches with FOS, Two CAWG cards in the core side, BMR embed into burst mode transceiver (BMT) card with FTL and the software OFP agent. The switching time of optical switches is 25 ns and the insert loss is lower than 4.5 dB. The frequency deviation of CAWG cards is 12.5GHz and the insert loss is 10.5 dB. The receiving power sensitivity of BMR is -25 dBm. The switching time of FTL has 98 ns and the frequency deviation is 2.5GHz according to ITU-T standard. The software OFP agent use PI to control the hardware through OFP. We use VMware software to build groups of virtual machines to realize Datacenters. Each virtual machine model a real node with the independent operation system, CPU and storage resource. In control plane, NC is realized with optical module control function, PCE computation function and resource abstraction function corresponding to three servers, database server is deployed to maintain transmission resources and the database of traffic engineering. AC server is used to carry TaSS strategy and monitoring the computing resources in servers. User plane is built in a server for running the required service.

We have designed experiment to verify the lightpath provisioning in CDON architecture for datacenter service migration. AC runs TaSS strategy to determine the path and schedule for service migration based on various application utilizations among datacenters and current network resource, then setup the path from source to destination node chosen by CSO. **Figure 7(b)-(d)** show the eye diagram and tuning waveform of FTL and spectrum of CAWG port reflected on the filter profile. The experimental results are further

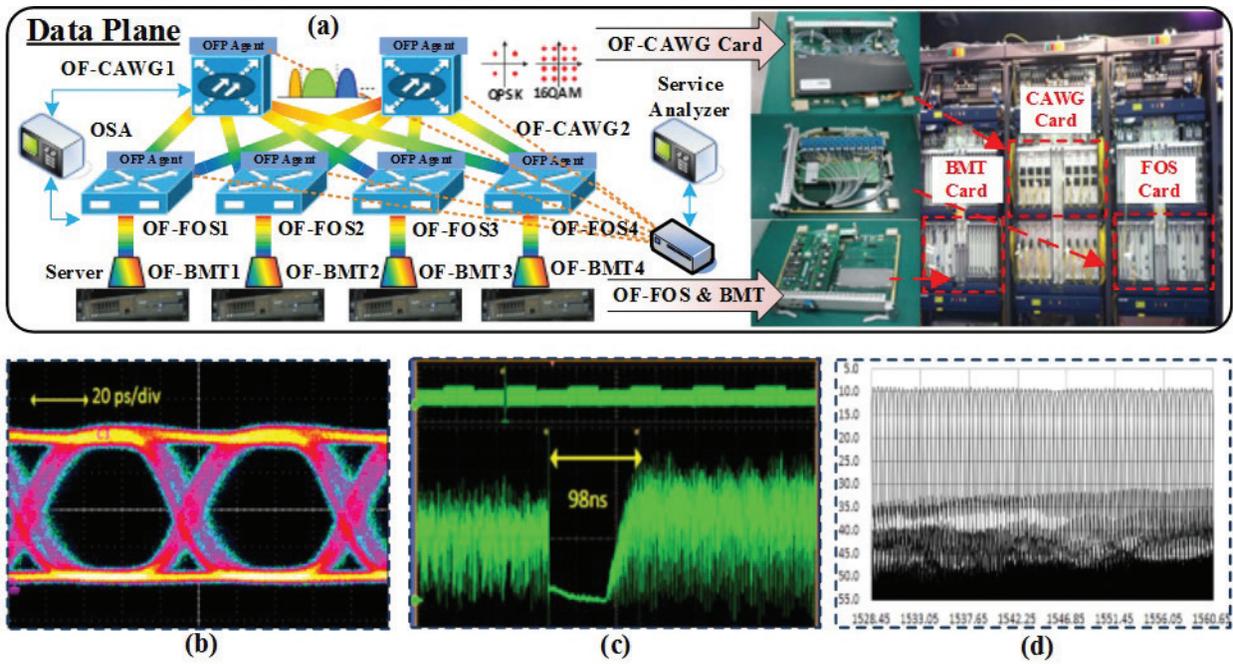


Figure 7. (a) Experimental testbed and demonstrator setup. (b, c) eye diagram and tuning waveform of FTL. (d) Spectrum of CAWG port.

shown in Figure 8(a) and (b). Figure 8(a) demonstrates the OpenFlow message exchange sequence for CDON by capturing in NC with Wireshark. In Figure 8, 10.108.50.74 denotes AC and 10.108.65.249 denotes NC, OF-OS nodes are 10.108.50.21 and 10.108.51.22. AC performs TaSS strategy and NC control all OF-OS via flow mod message through the allocated lighpath to provide the path. Then the path releasing information is sent via flow mod message, and keep synchronization by updating the application usage with UDP to. Figure 8(b) verifies OPF extensions for CDON in intra-datacenter networks with a snapshot of the extended flow table modification message for lighpath provisioning. We also evaluate the performance of CDON in intra-datacenter networks under the condition of heavy traffic load through simulation by comparing TaSS with traditional CSO (TCSO) strategy [16]. The migration data volumes from datacenter node following a Poisson process are randomly from 50Gbit to 400Gbit. Figure 9(a) and (b) compare the performances of two strategies. Comparing to TCSO, TaSS can reduce blocking probability effectively, especially when the network is under heavy load. We can also find that TaSS outperforms TCSO in the resource occupation rate significantly. The reason is TaSS realizes global optimization considering the time schedule with various delay sensitivity requirements and adjusts the service bandwidth according to the distribution of network resources.

in order to evaluate the proposed architecture in experiments, we build CDON including both control and data planes based on testbed, as shown in Figure 10(a). In data plane, SDN-based flexible optical nodes are deployed, each of which including flex ROADM and ODU boards.

Datacenters and the other nodes are also implemented on virtual machines created by VMware software on servers. As each virtual machine has independent operation system, IP address, CPU and memory resource, it can be regarded as a real node. The virtual OS technology can set up experiment topology including 14 nodes and 21 links which is similar to the backbone

of US. For OpenFlow-based CDON control plane, the NC is set up to support the proposed architecture and deployed in three servers corresponding to elastic spectrum control, physical layer parameter adjustment, PCE computation and resource abstraction, while the database server work for maintaining traffic engineering database, management information base, connection status and the configuration of the database and transport resources. The AC server support CSO agent and monitoring the application resources from datacenter networks with TaSS strategy. User plane is set up in a server for running the required application.

We have designed experiment to verify SDPs provisioning in CDON over flexi-grid optical networks for datacenter service migration whose results are shown in **Figure 10(b)-(d)**. AC runs TaSS strategy to determine the destination datacenter based on various application utilizations among datacenters and current network resource, and then set up SDP for the service migration along the lighpath from source to destination node. Moreover, for different SDP distances, the spectrum bandwidth and corresponding modulation format is tunable. **Figure 10(b)** and **(c)** show the spectrum of SDPs on the flexible link between two SD-OTNs which is reflected on the filter profile. The setup/release time of end-to-end SDP is evaluated and shown in **Figure 10(d)** by collecting the statistics of the strategy processing time, OFP transmission delay and device handle times of software and hardware.

Time.	Source	Destination	Protocol	Info
354.442690	10.108.50.74	10.108.65.249	UDP	Source port: 8
354.442797	10.108.65.249	10.108.50.74	UDP	Source port: 4
354.453564	10.108.50.74	10.108.65.249	UDP	Source port: 8
354.453633	10.108.65.249	10.108.50.21	OFP	Flow Mod (CSM)
354.453825	10.108.65.249	10.108.51.22	OFP	Flow Mod (CSM)
354.454894	10.108.50.21	10.108.65.249	OFP+Ethe	Packet In (AM)
354.454910	10.108.51.22	10.108.65.249	OFP+Ethe	Packet In (AM)
354.455391	10.108.65.249	10.108.50.74	UDP	Source port: 4
354.467208	10.108.65.249	10.108.50.21	OFP	Flow Mod (CSM)
354.467509	10.108.65.249	10.108.50.74	UDP	Source port: 3

(a)

```

    > Frame 2220 (172 bytes on wire, 172 bytes captured)
    > Linux cooked capture
    > Internet Protocol, Src: 10.108.65.249 (10.108.65.249), Dst: 10.108.50.21 (10.108.50.21)
    > Transmission Control Protocol, Src Port: 6633 (6633), Dst Port: 47232 (47232), Seq: 1041111111
    > OpenFlow Protocol
    > Header
    >   Version: 0x01
    >   Type: Flow Mod (CSM) (14)
    >   Length: 104
    >   Transaction ID: 0
    > Flow Modification
    0000 00 04 00 01 00 06 00 0c 29 7d 5b 2c 00 00 08 00  .....  )},.....
    0010 45 00 00 9c f6 be 40 00 40 06 ba b7 0a 6c 41 f9  E.....@ @....LA.
    0020 0a 6c 32 15 19 e9 b8 80 e7 70 89 cd a4 55 56 26  .l2.... .p...UV(
    0030 80 18 00 e3 49 e3 00 00 01 01 08 0a 00 92 00 0d  ..I...
    0040 09 29 3e 2e 01 0e 00 68 00 00 00 00 00 00 00  .>....h .....
    0050 0a 00 2b 00 74 76 e8 bf 7c 76 e8 bf 66 70 00 00  ..+.tv.. |v..hp..
    0060 e8 00 58 70 00 10 00 00 24 70 e8 bf e0 e4 2b 00  ..Xp.... $p....+.
    0070 e8 bf f4 df 10 00 00 02 00 00 00 01 00 00 00 06  .....
    0080 00 00 00 01 00 00 00 00 00 00 00 00 00 00 03 e7  .....
    0090 00 00 00 00 00 00 00 00 00 00 00 00 00 00 01  .....
    00a0 00 00 00 00 00 00 00 08 00 01 00 00
    
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(b)

Figure 8. (a)The capture of the OpenFlow message sequence and (b) extended flow mod message for CDON.

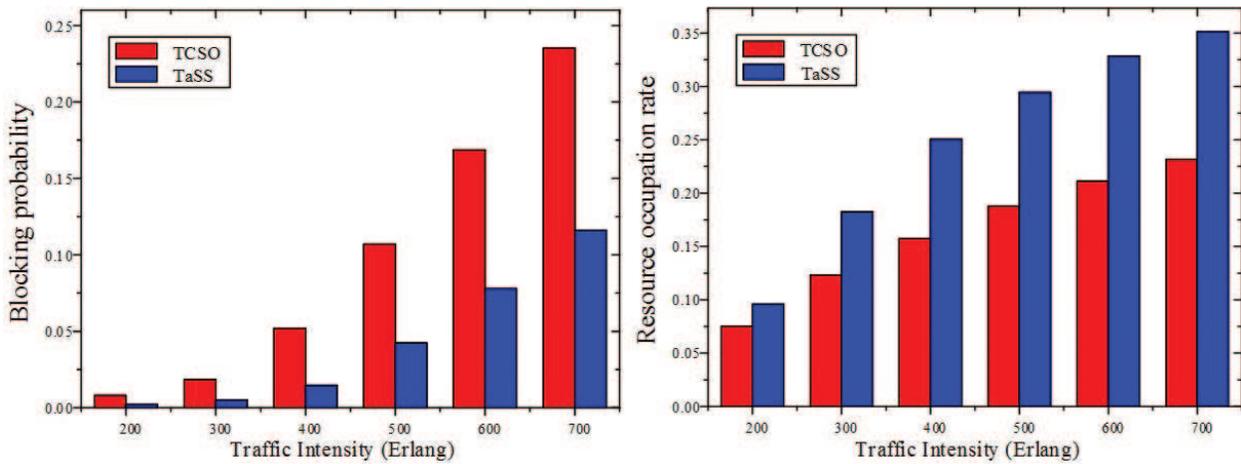


Figure 9. Comparisons on (a) blocking probability and (b) resource occupation rate between two strategies in heavy traffic load scenario.

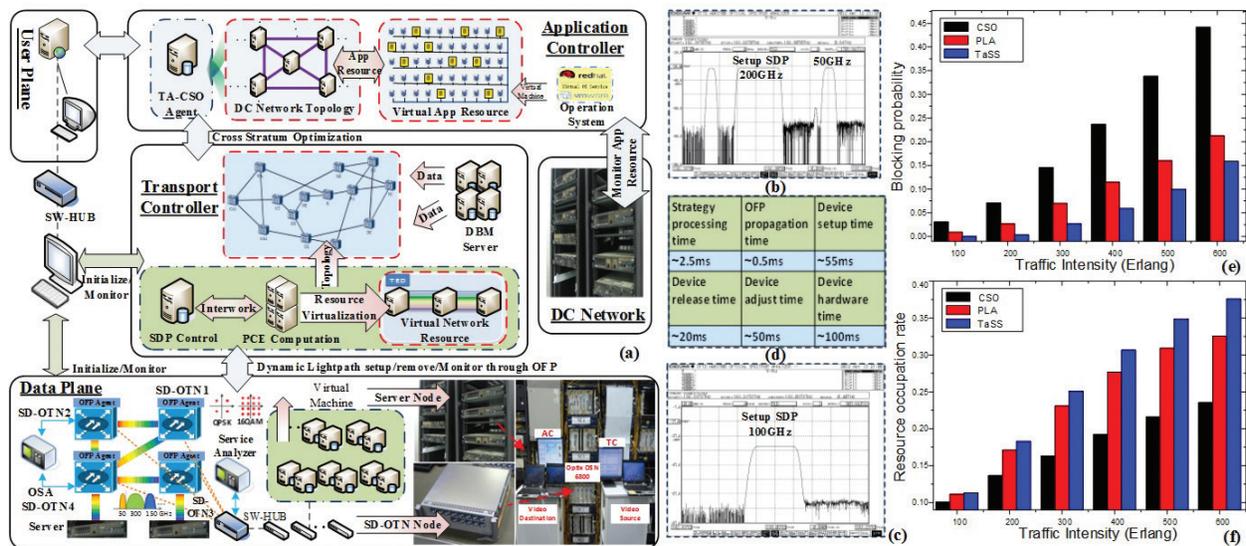


Figure 10. (a) Experimental testbed. (b, c) optical Spectrum of SDPs. (d) SDP setup/release delay. (e) Blocking probability (f) resource occupation rate under heavy load.

We test the performance of CDON by collecting the performance statistics in TaSS with CSO and physical layer adjustment strategies (PLA) under the condition of heavy load. The spectrum bandwidth of requests are randomly from 50GHz to 400GHz and the adjustable minimal frequency slot is 12.5GHz. All the destination nodes are datacentre in this network. These requests follow a Poisson process and statistics have been collected through the generation of 1×10^5 requests per execution. **Figure 10(e)** and **(f)** compare blocking probability and resource occupation rate of three strategies. In the figure, TaSS has a lower blocking probability than CSO and PLA, especially when the network is under the condition of heavy load. We can also find that TaSS can provide a better resource occupation rate than the other strategies. That is because both application and network resources can be considered in TaSS integrally and global optimization can be realized. Furthermore on the basis of it, by choosing high-level modulation format the spectrum resource is economized again. These results are further emphasized in **Figure 11(a)-(c)**.

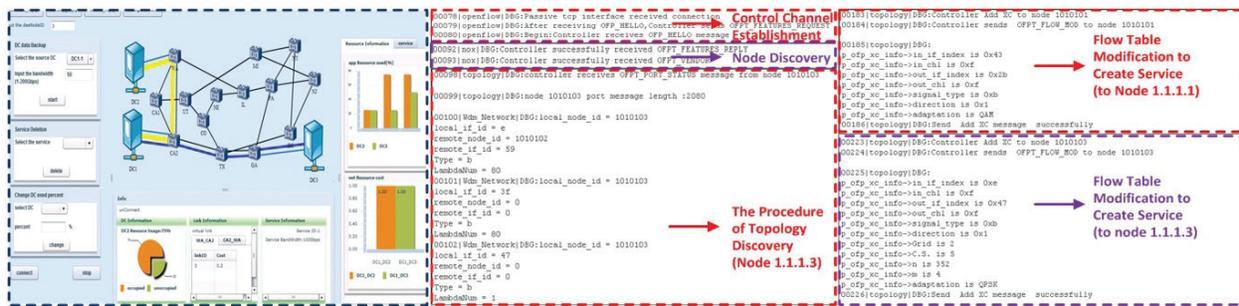


Figure 11. (a)Application graphical user interface (GUI) of testbed. (b) The capture of the OF messages for network discovery. (c) The capture of extended flow table message for SDP setup.

Figure 11(a) shows the application interface of CDON with status and destination node choice of datacenter service migration and various bandwidths of SDP. **Figure 11(b)** illustrates the OpenFlow message exchange for CDON with the capture in TC. **Figure 11(c)** shows a capture of the extended flow table modification message for SDP setup to verify the OPF extensions for CDON over flexi-grid optical networks.

5. Conclusion

Datacenter server resources will become more and more important in the future information society. It will be the main decision factor for reducing cost to allocate the resources efficiently. Considering the high capacity and low energy consumption requirements, Optical networking is regarded as promising technology for both intra-datacenter and inter-datacenter networks. The SDN technology can improve the resource utilization of application and network efficiently. In this chapter, we provide a CDON architecture in SDN-based datacenter optical network for service migration. Besides, we design the TaSS strategy for CDON and the extended OFP further. Based on our intra-datacenter and inter-datacenter networks testbed, we verify the feasibility and efficiency of CDON. The blocking probability and resource occupation rate of our approach under the condition of heavy traffic load are collected and compared with TaSS strategy. The experimental results indicate that the service with time sensitivity can be scheduled effectively and cross stratum resources utilization efficiency is improved in the CDON with TaSS strategy. Our future works for CDON are to improve TaSS performance with dynamic parameters and consider the scalability in optical network. Then we will realize the network virtualization in datacenter optical network on our OpenFlow-based testbed.

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