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Quick Response in a Continuous-Replenishment-Programme Based Manufacturer-Retailer Supply Chain

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1. Introduction

With the widespread concept of partnership in supply chain management, the traditional replenishment process is fast giving way to the Quick Response (QR) and Continuous Replenishment Programme (CRP). QR is a movement in industries to shorten the replenishment lead time which is critical to reduce inventory level and improve the levels of customer service. Wal-Mart, Seven-Eleven Japan, and many other retailers apply tremendous pressure on their suppliers to reduce the replenishment lead time (Chopra and Mendil 2007). CRP is an efficient replenishment initiative which focuses on removing excess inventory throughout the pipeline and synchronizes demand with production (EAN International 2000). In CRP, the inventory of retailer is planned, monitored, and replenished by the supplier on behalf of the consumers. To enable CRP, sales data and inventory level of the retailer must be provided to the supplier via Electronic Data Interchange (EDI) or other electronic means. Thus, to successfully implement CRP requires the supplier and the retailer to work in a cooperative manner based on mutual trust and joint gains.

The lead time is needed for several operations among trading parties, such as ordering, manufacturing, delivering, and handling. In practice, lead time can be shortened with extra crashing costs. Many researchers have studied the inventory decision incorporating lead time reduction under various assumptions. For example, Liao and Shyu (1991) first developed an inventory model in which lead time was the unique decision variable. Ben-Daya and Raouf (1994) extended the previous model by including both lead time and order quantity as decision variables. Pan and Yang (2002) and Pan and Hsiao (2005) considered lead time crashing cost as a function of both the order quantity and the reduced lead time. Since the reduction of lead time may involve with the vendor and the buyer to improve the related operations, a dyadic vendor-buyer viewpoint for shortening lead time is often suggested. For instance, Iyer and Bergen (1997) compare the profits of a single vendor and a single buyer supply chain before and after QR based on a newsboy inventory model. Ben-Daya and Hariga (2004), Ouyang et al. (2004), Chang et al. (2006), and Ouyang et al. (2007)

investigated the decision of lead time reduction in a context of a single-vendor-singleretailer integrated inventory system, in which the possibility of stockout of the supplier was ignored.

The study of inventory decision of CRP based inventory systems was initiated by Goyal (1976) by presenting a joint economic lot-size model for a supplier and a retailer. Subsequently, many researchers investigated this issue under various assumptions. For example, Banerjee (1986) generalized the model of Goyal (1976) by incorporating a finite replenishment rate for the supplier. Goyal (1988) extended Banerjee's work by relaxing the lot-for-lot policy for the supplier and assumed that the supplier's lot size is an integral multiple of retailer's order quantity. Goyal and Srinivasan (1992) further extended the model by relaxing the assumption that the supplier can supply to the retailer only after completing the entire production lot. At the same time, Banerjee and Banerjee (1992) extended the study of integrated inventory control to a multiple-retailer case by considering that the supplier delivers items to several buyers at a coordinated common cycle time. Banerjee and Banerjee (1994) generalized their previous model by dealing with the normally distributed demand case and assuming a fixed shortage cost attributable to each stockout incident.

To understand the effects of QR on the CRP based supply chains, we build a formal model of lead time reduction and replenishment decisions for a supply chain consisting of a manufacturer and multiple retailers in which the inventory throughout the supply chain is managed in a CRP context. The model extends the work of Banerjee and Banerjee (1994) by including the ordering cost and adopting QR with the expenditure regarded as a function of the reduced lead time. In stead of explicitly estimating the shortage cost, a service level constraint (SLC) approach is applied to the problem for which may skirt the difficult practical issue of explicitly determining the shortage cost (Chen and Krass 2001). According to the definition in Ouyang and Chuang (2000), the service level in this study is measured by the expected fraction of demand met from inventory on hand in an inventory cycle. The objective of the model is to determine the common shipment cycle time, and the target levels of replenishment for the manufacturer and each retailer so that the expected total system cost can be minimized under the service level constraint.

2. Notations and assumptions

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To develop the proposed models, we adopt the following notation which is principally the same as that in Banerjee and Banerjee (1994):

n = total number of retailers;

 d_i = demand rate of retailer *i*, which follows a probability density function (p.d.f.) $f_i(d_i)$ with mean D_i and variance σ_i^2 , i = 1, 2, ..., n;

 h_i = carrying cost per unit per year for retailer *i* (\$/unit/year);

 h_v = carrying cost per unit per year for the manufacturer (\$/unit/year);

- *l* = length of lead time for retailers (year), a decision variable;
- x_i = demand during the protection period (*T*+*l*) on retailer *i*, which has a p.d.f. $f_{i^{T+L}}(x_i)$ with the mean $D_i(T+l)$ and the variance $\sigma_i^2(T+l)$;

- *y* = demand during the production cycle time of *KT* on manufacturer, which has a p.d.f. *g*(*y*) with the mean $KT\sum_{i=1}^{n} D_i$ and the variance $KT\sum_{i=1}^{n} \sigma_i^2$;
- z_i = safety factor for retailer *i* with $z_i \ge 0$;
- z_v = safety factor for the manufacturer with $z_v \ge 0$;
- a_i = the threshold of retailer *i*'s service level, $0 \le a_i \le 1$;
- a_v = the threshold of manufacturer's service level, $0 \le a_v \le 1$;
- *A* = setup cost for the manufacturer (\$/per setup);
 - = common ordering cost shared by all retailers (\$/per order);
- C_i = individual ordering cost for retailer *i* (\$/per order);
- $C_r(l)$ = lead-time crashing cost per order;

P = production rate of manufacturer, which is a known constant;

D

С

= $\sum_{i=1}^{n} d_i$, total demand from all the retailers on the manufacturer per year, which

has mean
$$\sum_{i=1}^{n} D_i$$
 and variance $\sum_{i=1}^{n} \sigma_i^2$. Note that, of necessity, $P > \sum_{i=1}^{n} D_i$;

K = integral number of shipments to retailers per manufacturer's production cycle, which is a decision variable with $K \ge 1$;

- S_i = replenish-up-to level after placing a new order for retailer *i*, a decision variable;
- S_v = produce-up-to level after placing a new manufacture order for the manufacturer, a decision variable;

T = common shipment cycle time to all retailers (year), a decision variable;

 EC_i = the expected annual inventory cost for retailer *i*;

 EC_v = the expected annual inventory cost for the manufacturer;

ETC = the expected annual total cost for the system.

The major assumptions made in this study are the following:

- 1. In the CRP context, the manufacturer plans, monitors, and replenishes the inventory of retailers based on the information of sales data and inventory level provided by all the retailers to minimize the expected annual total cost of the supply chain.
- 2. The demand rate of each retailer is independently distributed.
- 3. Each production lot of the manufacturer will be delivered in an integral number, *K*, of shipments to all retailers periodically.
- 4. The target level of replenishment of S_i units for retailer *i* is equal to the sum of retailer *i*'s expected demand during the protection period, $D_i(T+l)$, and the safety stock, where safety stock = $z_i \sigma_i \sqrt{T+l}$.
- 5. The target level of production of S_v units for the manufacturer is equal to the sum of manufacturer's expected demand during the production cycle $KT\sum_{n=1}^{n}D_{n}$ and safety

manufacturer's expected demand during the production cycle, $KT\sum_{i=1}^{n}D_i$, and safety

stock, where safety stock = $z_v \sqrt{KT \sum_{i=1}^n \sigma_i^2}$.

- 6. For each order, a common ordering cost *C* is incurred to and shared by all retailers and an individual ordering cost C_i is incurred to the retailer *i*.
- 7. The manufacturer incurs a setup cost *A* for each production run.

- 8. The service level measures, i.e., fraction of demand met per cycle, for the retailer *i* and for the manufacturer are defined as $1 \frac{E(x_i S_i)^+}{D_i(T + l)}$ and $1 \frac{E(y S_v)^+}{KT\sum_{i=1}^n D_i}$, respectively.
- 9. The expenditure for implementing QR is modeled as the lead-time crashing cost per order, $C_r(l)$, and is assumed to be a non-increasing stairstep function of *l*, such as:

$$C_{r}(l) = \begin{cases} 0 & l = l_{0}, \\ r_{1} & l_{1} \le l < l_{0}, \\ \vdots & \vdots \\ r_{b} & l_{b} \le l < l_{b-1}, \end{cases}$$

where l_0 and l_b represent the existing and the minimal length of lead times, respectively. This cost could be expenditure on equipment improvement, order expediting, or special shipping and handling.

3. Model development and solution

Figure 1 depicting the inventory time plots for a manufacturer and *n* retailers, illustrates the CRP based inventory control system proposed in this work. For the retailer *i*, there is a target level of replenishment of S_i units. For the manufacturer, there is a target level of production of S_v units. Given a common shipment cycle time of *T* years, the manufacturer produces a batch of item every *KT* years. After the *K*th shipment cycle since the beginning of the last production cycle, the manufacturer begins the production of another lot DT/P years prior to the next scheduled shipment time, as shown in Figure 1, which is illustrated the case of *K* =3. In order to cope with a stochastic demand, all of the trading parties carry safety stocks.

The demand during the protection period for retailer *i* is assumed to be normally distributed with mean $D_i(T+l)$ and standard deviation $\sigma_i\sqrt{T+l}$. Then, the total demand from all retailers on the manufacturer during a production cycle of *KT* is normally distributed with

the mean $KT\sum_{i=1}^{n} D_i$ and the standard deviation $\sqrt{KT\sum_{i=1}^{n} \sigma_i^2}$. For the retailer *i*, with a cycle of

T, a replenishment lead time of *l*, and a target level of S_i , the expected net inventory level just before receipt of an order is given by S_i - $D_i l$ - $D_i T$ and the expected net inventory level immediately after the successive order is S_i - $D_i l$. Then, the average inventory over the cycle is

approximated by $\frac{1}{2} \left[(S_i - D_i l - D_i T) + (S_i - D_i l) \right]$. According to the fourth assumption, the target levels of replenishments, S_i , is set as

$$S_i = D_i \left(T + l \right) + z_i \sigma_i \sqrt{T + l} . \tag{1}$$

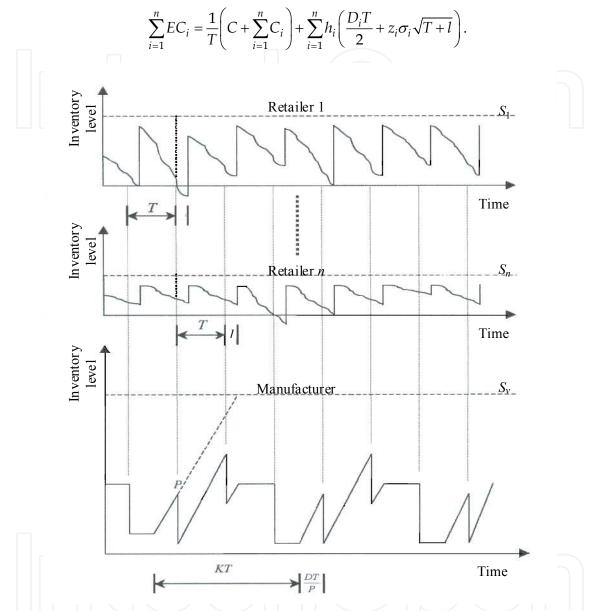
So that the retailer *i*'s expected annual carrying cost is

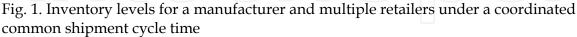
$$\frac{h_i}{2}\Big[\big(S_i - D_i l - D_i T\big) + \big(S_i - D_i l\big)\Big] = h_i \bigg(S_i - D_i l - \frac{D_i T}{2}\bigg) = h_i \bigg(\frac{D_i T}{2} + z_i \sigma_i \sqrt{T + l}\bigg).$$

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The annual individual ordering cost incurred to retailer *i* is C_i/T . Then, the expected annual inventory cost for all retailers, consisting of common ordering cost, individual ordering cost, and carrying cost, is given by





Based on the fifth assumption, the target levels of replenishments, S_v , is set as

$$S_{v} = KT \sum_{i=1}^{n} D_{i} + z_{v} \sqrt{KT \sum_{i=1}^{n} \sigma_{i}^{2}} .$$
⁽²⁾

According to Banerjee and Banerjee (1994) with Eq. (2), the manufacturer's average inventory can be derived as

$$\frac{\sum_{i=1}^{n} D_{i}T}{2} \left[\frac{\sum_{i=1}^{n} D_{i}}{P} (2-K) + K - 1 \right] + S_{v} - KT \sum_{i=1}^{n} D_{i} = \frac{\sum_{i=1}^{n} D_{i}T}{2} \left[K \left(1 - \frac{\sum_{i=1}^{n} D_{i}}{P} \right) + \frac{2\sum_{i=1}^{n} D_{i}}{P} - 1 \right] + z_{v} \sqrt{KT \sum_{i=1}^{n} \sigma_{i}^{2}}.$$

Then, the manufacturer's expected annual inventory cost is given by

$$EC_{v} = \frac{A}{KT} + h_{v} \left\{ \frac{\sum_{i=1}^{n} D_{i}T}{2} \left[K \left(\frac{\sum_{i=1}^{n} D_{i}}{1 - \frac{i=1}{P}} \right) + \frac{2\sum_{i=1}^{n} D_{i}}{P} - 1 \right] + z_{v} \sqrt{KT} \sum_{i=1}^{n} \sigma_{i}^{2} \right\},$$

where the first term represents the average annual setup cost and the second term is the average carrying cost.

Based on the service level measure defined in Ouyang and Chuang (2000), the following service level constraints are specified:

service level constraint on retailer
$$i = 1 - \frac{\int_{S_i}^{\infty} (x_i - S_i) f_{i-T+l}(x_i) dx_i}{D_i(T+l)} \ge a_i, i = 1, 2, \dots, n,$$
 (3)

service level constraint on the manufacturer = $1 - \frac{\int_{S_v}^{\infty} (y - S_v) g(y) dy}{KT \sum_{i=1}^{n} D_i} \ge a_v.$ (4)

As the demand is normally distributed, (3) and (4) can be rewritten as (see Appendix for proof):

service level constraint on retailer
$$i = \phi(z_i) - z_i \Phi(z_i) - \frac{(1 - \alpha_i)D_i\sqrt{T + l}}{\sigma_i} \le 0$$
, $i = 1, 2, ..., n$, (5)
service level constraint on the manufacturer $= \phi(z_v) - z_v \Phi(z_v) - \frac{(1 - \alpha_v)\sqrt{KT}\sum_{i=1}^n D_i}{\sqrt{\sum_{i=1}^n \sigma_i^2}} \le 0$, (6)

where $\phi(z)$ and $\Phi(z)$ represent the standard normal distribution and the complementary cumulative normal distribution function, respectively. The expected annual total cost for the system, comprising the expected annual inventory costs for all parties and the lead time crashing cost, is given by

$$ETC(K, T, l, \underline{z}, z_{v}) = \sum_{i=1}^{n} EC_{i} + EC_{v} + \frac{C_{r}(l)}{T}$$

$$= \frac{C + \sum_{i=1}^{n} C_{i} + \frac{A}{K} + C_{r}(l)}{T} + T \left\{ \frac{h_{v} \sum_{i=1}^{n} D_{i}}{2} \left[K \left(1 - \frac{\sum_{i=1}^{n} D_{i}}{P} \right) + \frac{2\sum_{i=1}^{n} D_{i}}{P} - 1 \right] + \sum_{i=1}^{n} \frac{D_{i}h_{i}}{2} \right\}$$

$$+ h_{v} z_{v} \sqrt{KT \sum_{i=1}^{n} \sigma_{i}^{2}} + \sum_{i=1}^{n} h_{i} z_{i} \sigma_{i} \sqrt{T + l} , \qquad (7)$$

where the vector $\underline{z} = \{z_1, z_2, ..., z_n\}$. Then, the decision is to minimize $ETC(K, T, l, \underline{z}, z_v)$ under the constraints (5) and (6).

In order to solve the nonlinear programming problem, the following propositions are needed.

Proposition 1. For any given *K*, *T*, \underline{z} , and z_v , the minimal *ETC* will occur at the end points of the interval $[l_{j+1}, l_j]$ when $C_r(l)$ is a stairstep function. Proof:

As $C_r(l)$ is a stairstep function, $\frac{\partial^2 C_r(l)}{\partial l^2} = 0$ for any l in a continuous interval $[l_{j+1}, l_j]$. The second derivative of (7) with respect to l is equal to

$$\frac{\partial^2 ETC}{\partial l^2} = \frac{1}{T} \frac{\partial^2 C_r(l)}{\partial l^2} - \frac{\sum_{i=1}^n h_i z_i \sigma_i}{4\sqrt{(T+l)^3}} < 0$$

for any *l* in a continuous interval $[l_{j+1}, l_j]$. Therefore, the proposition holds. **Proposition 2.** For any fixed *K*, *l*, *<u>z</u>, and <i>z*_v, there is a unique optimal *T* to *ETC* in (7). Proof:

For any given *K*, *l*, \underline{z} , and z_v in a continuous interval $[l_{j+1}, l_j]$, a necessary condition for a solution of *T* to be optimal for (7) is

$$\frac{\partial ETC}{\partial T} = \frac{C + \sum_{i=1}^{n} C_i + \frac{A}{K} + C_r(l)}{-T^2} + \frac{h_v \sum_{i=1}^{n} D_i}{2} \left[K \left(1 - \frac{\sum_{i=1}^{n} D_i}{P} \right) + \frac{2\sum_{i=1}^{n} D_i}{P} - 1 \right] + \frac{\sum_{i=1}^{n} D_i h_i}{2} + \frac{h_v z_v \sqrt{K \sum_{i=1}^{n} \sigma_i^2}}{2\sqrt{T}} + \frac{\sum_{i=1}^{n} h_i z_i \sigma_i}{2\sqrt{T + l}} = 0.$$

The above equation can be rearranged as

$$\frac{C + \sum_{i=1}^{n} C_{i} + \frac{A}{K} + C_{r}(l)}{T^{2}} = \frac{h_{v} \sum_{i=1}^{n} D_{i}}{2} \left[K \left(1 - \frac{\sum_{i=1}^{n} D_{i}}{P} \right) + \frac{2\sum_{i=1}^{n} D_{i}}{P} - 1 \right] + \sum_{i=1}^{n} \frac{D_{i}h_{i}}{2} + \frac{h_{v} z_{v} \sqrt{K \sum_{i=1}^{n} \sigma_{i}^{2}}}{2\sqrt{T}} + \frac{\sum_{i=1}^{n} h_{i} z_{i} \sigma_{i}}{2\sqrt{T + l}} \right]$$
(8)

Let T_f be the solution to (8). Then, the second derivative of (7) at T_f satisfies

$$\begin{split} \frac{\partial^2 ETC}{\partial T^2(T_f)} &= \frac{2\left(C + \sum_{i=1}^n C_i + \frac{A}{K} + C_r(l)\right)}{T_f^3} - \frac{h_v z_v \sqrt{K \sum_{i=1}^n \sigma_i^2}}{4T_f^{\frac{3}{2}}} - \frac{\sum_{i=1}^n h_i z_i \sigma_i}{4\left(T_f + l\right)^{\frac{3}{2}}} \\ &= \frac{2}{T_f} \left\{ \frac{h_v \sum_{i=1}^n D_i}{2} \left[K \left(1 - \frac{\sum_{i=1}^n D_i}{P} \right) + \frac{2\sum_{i=1}^n D_i}{P} - 1 \right] + \sum_{i=1}^n \frac{D_i h_i}{2} + \frac{h_v z_v \sqrt{K \sum_{i=1}^n \sigma_i^2}}{2\sqrt{T_f}} + \frac{\sum_{i=1}^n h_i z_i \sigma_i}{2\sqrt{T_f + l}} \right\} \\ &- \frac{h_v z_v \sqrt{K \sum_{i=1}^n \sigma_i^2}}{4T_f^{\frac{3}{2}}} - \frac{\sum_{i=1}^n h_i z_i \sigma_i}{4\left(T_f + l\right)^{\frac{3}{2}}} \\ &= \frac{h_v \sum_{i=1}^n D_i}{T_f} \left[K \left(1 - \frac{\sum_{i=1}^n D_i}{P} \right) + \frac{2\sum_{i=1}^n D_i}{P} - 1 \right] + \sum_{i=1}^n \frac{D_i h_i}{T_f} + \frac{3h_v z_v \sqrt{K \sum_{i=1}^n \sigma_i^2}}{4T_f \sqrt{T_f}} + \frac{\left(3T_f + 4l\right) \sum_{i=1}^n h_i z_i \sigma_i}{T_f \left(T_f + l\right)^{\frac{3}{2}}} > 0. \end{split}$$

That is, the second derivative of *ETC* at each T_f is positive. The result implies that, for given *K*, *l*, *z*, and *z*_v, an only solution for *T*, corresponding to the minimum of *ETC*, can be derived from (8). This completes the proof.

Proposition 3. For any given *K*, *T*, and *l* in a continuous interval $[l_{j+1}, l_j]$, the boundary solution for \underline{z} and z_v derived from (5) and (6) is the optimal solution to this nonlinear programming problem.

Proof:

By using the method of Lagrange multipliers, the problem can be transformed into

$$ETC(K, T, l, \underline{z}, z_{v}, \underline{\lambda}, \underline{s}) = \frac{C + \sum_{i=1}^{n} C_{i} + \frac{A}{K} + C_{r}(l)}{T} + T \left\{ \frac{h_{v} \sum_{i=1}^{n} D_{i}}{2} \left[K \left(\sum_{i=1}^{n} D_{i} \right) + \frac{2\sum_{i=1}^{n} D_{i}}{P} - 1 \right] + \sum_{i=1}^{n} \frac{D_{i}h_{i}}{2} \right\} + h_{v} z_{v} \sqrt{KT} \sum_{i=1}^{n} \sigma_{i}^{2} + \sum_{i=1}^{n} h_{i} z_{i} \sigma_{i} \sqrt{T + l} + \sum_{i=1}^{n} \lambda_{i} \left[\phi(z_{i}) - z_{i} \Phi(z_{i}) - \frac{(1 - \alpha_{i})D_{i} \sqrt{T + l}}{\sigma_{i}} + s_{i}^{2} \right] + \lambda_{v} \left[\phi(z_{v}) - z_{v} \Phi(z_{v}) - \frac{(1 - \alpha_{v})\sqrt{KT} \sum_{i=1}^{n} D_{i}}{\sqrt{\sum_{i=1}^{n} \sigma_{i}^{2}}} + s_{v}^{2} \right], \qquad (9)$$

where $\underline{\lambda}$ is a vector of Lagrange multipliers with $\underline{\lambda} = \{\lambda_1, \lambda_2, ..., \lambda_n, \lambda_v\} \ge 0$, and \underline{s} is a vector of slack variables with $\underline{s} = \{s_1^2, s_2^2, ..., s_n^2, s_v^2\}$. According to Kuhn-Tucker theorem (Taha, 2002), for any given l in a continuous interval $[l_{j+1}, l_j]$, it can be shown that $\underline{s} = \{0, 0, ..., 0\}$ is a necessary condition for a solution in (9) to be optimal. Then, for any given l in a continuous interval $[l_{j+1}, l_j]$, the necessary conditions for (9) to be minimized are

$$\frac{\partial ETC}{\partial z_i} = h_i \sigma_i \sqrt{T+l} - \lambda_i \Phi(z_i) = 0 , i = 1, 2, ..., n,$$
(10)

$$\frac{\partial ETC}{\partial z_v} = h_v \sqrt{KT \sum_{i=1}^n \sigma_i^2} - \lambda_v \Phi(z_v) = 0 , \qquad (11)$$

$$\frac{\partial ETC}{\partial \lambda_i} = \phi(z_i) - z_i \Phi(z_i) - (1 - \alpha_i) D_i \frac{\sqrt{T+l}}{\sigma_i} = 0, \qquad (12)$$

$$\frac{\partial ETC}{\lambda_v} = \phi(z_v) - z_v \Phi(z_v) - \frac{(1 - \alpha_v)\sqrt{KT} \sum_{i=1}^n D_i}{\sqrt{\sum_{i=1}^n \sigma_i^2}} = 0, \qquad (13)$$

From (10) and (11), we have

$$\lambda_{i}^{*}(z_{i}) = \frac{h_{i}\sigma_{i}\sqrt{T+l}}{\Phi(z_{i})} > 0, \ i = 1, 2, ..., n, \text{ and } \lambda_{v}^{*}(z_{v}) = \frac{h_{v}\sqrt{KT\sum_{i=1}^{n}\sigma_{i}^{2}}}{\Phi(z_{v})} > 0.$$

Therefore, constraints (5) and (6) are binding. In addition, $\frac{\partial^2 ETC}{\partial z_i^2} = \lambda_i \phi(z_i) > 0$ (*i*=1, 2,..., *n*,

m), then the \underline{z} and z_v derived from (12) and (13) (i.e., the boundary solutions to (5) and (6)) are optimal to (9) with given *K*, *T*, and *l*.

Based on the above propositions, the optimal solution for *K*, *T*, *l*, \underline{S} , and S_v to this nonlinear programming problem can be obtained by the following algorithm:

- **Algorithm** 10. Set *K*=1.
- 11. Find $ETC^*(K, l_i | j = 0, 1, ..., b)$
 - 2.1 Let *j* =0
 - 2.2 Set $l = l_j$. z_i (i = 1, 2, ..., n) = 0, and $z_v = 0$.
 - 2.3 Compute T^* by substituting K, l_j , \underline{z} and z_v into (8). Set $T_{kj} = T^*$.
 - 2.4 Compute \underline{z}^* by substituting l_j and T_{kj} into (12) and compute z_v^* by substituting *K* and T_{kj} into (13).
 - 2.5 Compute *T*^{*} by substituting *K*, l_j , \underline{z}^* , and z_v^* into (8).
 - 2.6 If $T^* = T_{kj}$, go to step 2.7; otherwise, set $T_{kj} = T^*$ and go to step 2.4.
 - 2.7 Set $T^*(K,l_j) = T_{kj}$, $\underline{z}^*(K,l_j) = \underline{z}^*$, and $z_v^*(K,l_j) = z_v^*$. Compute $ETC^*(K,l_j)$ by substituting

K, *l_j*, *T*^{*}(*K*,*l_j*), $\underline{z}^{*}(K, l_{j})$, and $z_{v}^{*}(K, l_{j})$ into (7).

2.8 Let j = j+1 and go back to step 2.2 until j = b.

- 12. Set $ETC(K,l) = Min[ETC(K,l_j) | j = 0,1,...,b].$
- 13. If K = 1, set $ETC_s = ETC(K,l)$, K = K+1, z_i (i = 1, 2, ..., n) = 0, $z_v = 0$ and go back to step 2; otherwise go to step 5.
- 14. If $ETC(K,l) < ETC_s$, let $ETC_s = ETC(K,l)$, K = K+1, z_i (i = 1, 2, ..., n) = 0, $z_v = 0$, and go back to step 2; otherwise let $ETC^* = ETC_s$ along with the corresponding values of K and l being the value of K^* and l^* , respectively.
- 15. Find T^* from (8) with the obtained K^* and l^* . Find \underline{z}^* from (12) and z_v^* from (13) with the obtained K^* , l^* , and T^* . Find \underline{S}^* from (1) and S_v^* from (2) with the obtained K^* , l^* , T^* , \underline{z}^* , and z_v^* . Output the values of K^* , l^* , T^* , \underline{S}^* , S_v^* , ETC^* and Stop.

4. Computer implementation

The algorithm proposed in Section 3 has been implemented as a decision support system on a personal computer in which Intel Pentium D 2.8 GHz CPU with 1024 MB RAM inside it. Visual Basic 2005 is utilized as the software platform to develop the decision support system. Figure 2 shows the window to input the numbers of retailers, e.g., n = 3 in this case, and steps of the lead-time crashing cost function, e.g., three steps in this case, involved in the replenishment and lead time reduction decisions. Once the decision maker inputs the numbers, a window to specify the relevant parameters will be displayed as shown in Figure 2. If all of the relevant parameters are input for the system, the decision maker then clicks the label "Find Solution" to select the lead time policy and to set the options of output. Figure 3 illustrates the result that the system derives the solution in which $K^* = 2$, $l^* = 0.005$ years, $T^* = 0.0709$ years, $S_1^* = 708$, $S_2^* = 760$, $S_3^* = 1276$, $S_v^* = 3574$, and $ETC^* = \$19455.5$ for the case of implementing QR (representing by the controllable lead time option) after the

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	rameters		olution	What-i	Vhat-if Analysis		
nput the number nput the numbe		 lead-time crashing cost	function: <u>3</u>	Enter			
nput parameter	s:						
Manufacturer:							
Production rate	Setup cost	Carrying cost rate	Minimal service level				
28000	200	3	0.99				
Retailers	Demand rate	Standard deviation of demand rate	Minor order cost	Carrying cost rate	Minimal service level		
Retailers l	6000	600	100	5	0.99		
Retailers2	5000	800	150	4	0.99		
Retailers3	10000	900	80	4,5	0.99		

Fig. 2. The system window to input parameters

Input Parameters	Find Solution	What-if Analysis				
Select Lead Time Policy:	Output the Solution	n:				
Controllable lead time	🔽 Number of Shipm	ents per Production Cycle				
Constant lead time	🔽 Lead Time					
	🔽 Shipments Cycle Time					
	🔽 Replenish-up-to I	evel				
	🔽 Total System Cos	t				
	🗖 Total Cost of Ma	mufacturer				
	🗖 Total Cost of Ret	tailers				
	🔽 Benefit of Lead-t	une Perform				
The Solution K*=2 1*=0.005 T*=0.0709 S1*=708 S2*=760 S3*=1276 Sy *=3574						
$ETC^* = 19455.5$						

Fig. 3. The replenishment and lead time reduction decisions under CRP provided by the decision support system

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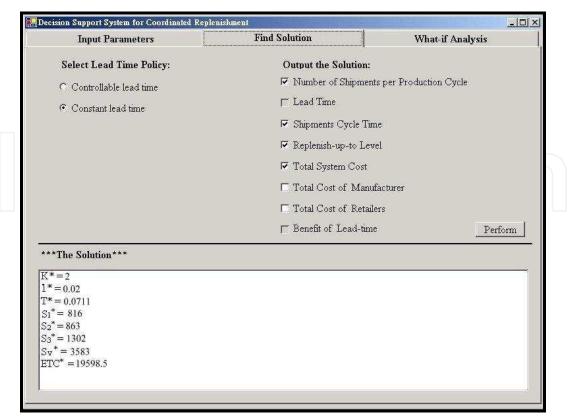


Fig. 4. The coordinated replenishment decision under CRP provided by the decision support system

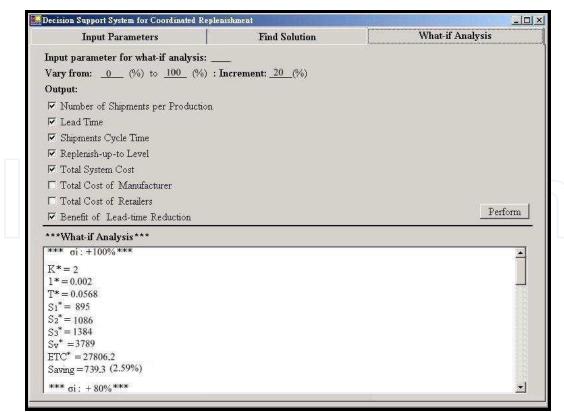


Fig. 5. The "what-if analysis" function of the decision support system

decision maker clicks "Perform". From the value of "Saving", the benefit derived from implementing QR can be learned. As shown in Figure 4, the system also provides the inventory decision before QR (represented by the constant lead time option). One advantage of the system is that the decision maker can learn the sensitivity of replenishment and lead time decisions and the expected annual total cost of the chain by specifying the parameter with the variation range in the Window of "What-if Analysis". Figure 5 illustrates the output of the what-if analysis of K^* , l^* , T^* , \underline{S}^* , S_v^* , *ETC**, and the cost saving after QR against the variation of standard deviation of demand.

		K^*	T^*	<i>l</i> *	S_1^*	S_2^*	S_3^*	S_v^*	ETC*	K_0^a	T_0^{a}	ΔETC(%) ^ь
Base ex	ample	2	0.0709	0.005	708	760	1276	3574	19455.3	2	0.0711	0.73
σ _i (i=1,2,3)	+100%	2	0.0568	0.002	895	1086	1384	3789	27806.2	2	0.0577	2.59
Р	+100%	1	0.0877	0.005	827	874	1320	2357	18822.3	1	0.0880	0.71
С	+100%	2	0.0779	0.005	758	808	1208	3885	20799.7	2	0.0781	0.67
C_i (<i>i</i> =1,2,3)	+100%	1	0.1024	0.005	930	970	1488	2691	23410.0	1	0.1027	0.54
Α	+100%	3	0.0705	0.005	705	757	1291	5108	20696.8	2	0.0781	1.16
h_i (<i>i</i> =1,2,3)	+100%	3	0.0507	0.005	560	616	886	3803	26567.3	3	0.0513	1.88
h_v	+100%	1	0.0704	0.005	704	756	1120	1958	23078.6	1	0.0705	0.62
$1-a_i$ (<i>i</i> =1,2,3)	+100%	2	0.0733	0.01	705	741	1128	3681	18548.2	2	0.0735	0.32
$1-a_v$	+100%	2	0.0727	0.005	721	772	1147	3467	18897.2	2	0.0728	0.75

a: K_0 and T_0 represent the optimal K and T, respectively, before QR is implemented.

b: $\Delta ETC(\%) = (1 - ETC^* / ETC_0) \times 100\%$, where ETC_0 represents the optimal ETC before QR is implemented.

 Table 1. Sensitivity analysis for parameters

The sensitivity of the optimal solution has been further examined by conducting several numerical experiments. The results are illustrated by a base case and are shown in Table 1. Some findings are summarized as follows:

- 1. The amount of $\Delta ETC(\%)$ measures the improvement rate of expected annual total system cost after implementing QR. The behavior of $\Delta ETC(\%)$ reveals that the supply chain may benefit from shortening the replenishment lead time.
- 2. The value of $\Delta ETC(\%)$ is especially sensitive to the variation of σ_i or h_i . As σ_i or h_i increases, the value of $\Delta ETC(\%)$ increases. The results mean that the return of investment in QR is especially significant for a system with high uncertainty in demand or with high carrying costs incurred to the retailers.
- 3. The common shipment cycle time is more sensitive to the variance of demand, the variation of manufacturer's production rate, the individual ordering cost, and the retailer's carrying cost.
- 4. The shipment cycle time, T^* , will be reduced after implementing QR in which the replenishment lead time for retailers has been shortened. Moreover, the target levels of

replenishments for the manufacturer and the retailers may become lower after bringing QR into practice.

- 5. The number of shipments per production cycle, *K**, after implementing QR is always no less than that before implementing QR. When the values of *K* for the two situations are equal, the optimal shipment cycle time, *T** as well as the protection period, *T**+*l**, after QR is always no large than those before QR. The result implies that under a fixed number of shipments during a production cycle, the protection period for the retailer will be reduced after implementing QR.
- 6. The increase of production rate will result in a longer shipment cycle time and higher target levels of replenishments for the retailers. In contrast, as the value of *P* increases, the number of shipment per production cycle and the target level of production for the manufacturer will become smaller.
- 7. When the ordering cost increases, the shipment cycle time and the target levels of replenishments for each party will increase.
- 8. The value of a_i specifies the minimal fraction of demand met per cycle for the retailer and directly relates with the length of protection period (*T*+*l*). It can be found that the amount of $\Delta ETC(\%)$ decreases as the retailer's maximal fraction of demand unfilled per cycle increases. The result implies that the benefit from implementing QR is significantly related to the retailer's service level threshold and the benefit is substantial for a supply chain requesting a high customer service level.
- 9. The numerical example shows that as $1-a_v$ increases from 1% to 2%, the length of common shipment cycle time increases by 2.54% (from 0.0709 to 0.0727) but the lead time is unaffected. The result implies that the effect of manufacturer's service level threshold is more significant on the common shipment cycle time than on the reduction of lead time.

5. Conclusions

Speed, service, and supply chain management have been core capabilities for business competition. The reduction of overall system response time with a satisfied service level has received a great deal of attention from researchers and practitioners. In this study, we investigated the effect of investing in QR on a CRP based supply chain where a manufacturer produces and delivers items to multiple retailers at a coordinated common cycle time with the minimized expected total system cost and satisfied service levels. Extending the work of Banerjee and Banerjee (1994) by involving ordering costs and the reducible replenishment lead time, a model and an algorithm are proposed to simultaneously determine the optimal shipment cycle time, target levels of replenishments, lead time, and number of shipments pre production cycle under service level constraints for the supply chain.

A numerical experiment along with sensitivity analysis was performed and the results explain the effect of QR on the replenishment decisions and the total system cost. The results provide the following findings about our model:

- 1. The system can reduce the stocks throughout the pipeline and remain service levels of retailers via investing in QR initiative.
- 2. The benefit from implementing QR is especially significant for a supply chain with high uncertainty in demand or the retailers requesting high service levels or incurring high carrying costs.

- 3. The shipment cycle time will decrease after QR implemented. Additionally, the shipment cycle time is especially sensitive to the variation of manufacturer's production rate, the individual ordering cost, the variance of demand, and the retailer's carrying cost.
- 4. The decision of adopting QR is mainly influenced by the variance of demand and the retailer's service level threshold. The higher the demand uncertainty or the higher the retailer's service level threshold, the more beneficial to implement QR in supply chains.

Appendix

Let $G = \int_{S_i}^{\infty} (x - S_i) f(x_i) dx_i$ and $z_i = \frac{x_i - \mu_i}{\theta_i}$, where $f(x_i)$ represents a normal distribution with

the mean μ_i and the standard deviation θ_i . By replacing x_i with $\mu_i + z_i \theta_i$, we have

$$G = \int_{\underline{S_i - \mu_i}}^{\infty} \left(\mu_i + z_i \theta_i - S_i \right) \phi(z_i) dz_i = \sigma_i \int_{\underline{S_i - \mu_i}}^{\infty} z_i \phi(z_i) dz_i - \left(S_i - \mu_i \right) \int_{\underline{S_i - \mu_i}}^{\infty} \phi(z_i) dz_i , \qquad (A.1)$$

where $\phi(z_i)$ is the standard normal distribution and $\Phi(z_i)$ is the complementary cumulative normal distribution function defined as

$$\Phi(z_i) = \int_{z_i}^{\infty} \phi(w) dw = \frac{1}{\sqrt{2\pi}} \int_{z_i}^{\infty} e^{-\frac{w^2}{2}} dw$$

Let $v = \frac{z_i^2}{2}$, then there is $dv = z_i dz_i$ and

$$\int_{\frac{S_{i}-\mu_{i}}{\theta_{i}}}^{\infty} z_{i}\phi(z_{i})dz = \frac{1}{\sqrt{2\pi}} \int_{\frac{1}{2} \left[\frac{(S_{i}-\mu_{i})}{\theta_{i}}\right]^{2}}^{\infty} e^{-v}dv = \frac{-1}{\sqrt{2\pi}} e^{-v} \left| \frac{1}{\frac{1}{2} \left[\frac{(S_{i}-\mu_{i})}{\theta_{i}}\right]^{2}} = \frac{1}{\sqrt{2\pi}} e^{-\frac{1}{2} \left(\frac{S_{i}-\mu_{i}}{\theta_{i}}\right)^{2}} = \phi(z_{i}) \text{ (A.2)}$$

Next,

$$\left(S_{i}-\mu_{i}\right)\int_{\underline{S_{i}-\mu_{i}}}^{\infty}\phi(z_{i})dz_{i}=\left(S_{i}-\mu_{i}\right)\Phi(z_{i}).$$
(A.3)

By substituting (A.2) and (A.3) into (A.1), we have

$$G = \int_{S_i}^{\infty} (x - S_i) f(x_i) dx_i = \theta_i \phi(z_i) - (S_i - \mu_i) \Phi(z_i) .$$
(A.4)

Let $\mu_i = D_i(T+l)$ and $\theta_i = \sigma_i \sqrt{T+l}$ and apply (A.4) to (3), (3) can be rewritten as

$$1 - \frac{\sigma_i \sqrt{T + l}\phi(z_i) - \left[S_i - D_i \left(T + l\right)\right] \Phi(z_i)}{D_i \left(T + l\right)} \ge a_i, \ i = 1, 2, 3.$$
(A.5)

By substituting (1) in to (A.5) and rearranging the inequality, we have

$$\phi(z_i) - z_i \Phi(z_i) - (1 - \alpha_i) D_i \frac{\sqrt{T+l}}{\sigma_i} \le 0, i = 1, 2, 3.$$

This completes the proof for (5). Accordingly, (6) can be proved.

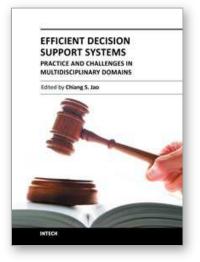
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